

The CEO Problem With Secrecy Constraints

Farshad Naghibi, *Student Member, IEEE*, Somayeh Salimi, *Member, IEEE*,
and Mikael Skoglund, *Senior Member, IEEE*

Abstract—We study a lossy source coding problem with secrecy constraints in which a remote information source should be transmitted to a single destination via multiple agents in the presence of a passive eavesdropper. The agents observe noisy versions of the source and independently encode and transmit their observations to the destination via noiseless rate-limited links. The destination should estimate the remote source based on the information received from the agents within a certain mean distortion threshold. The eavesdropper, with access to side information correlated to the source, is able to listen in on one of the links from the agents to the destination in order to obtain as much information as possible about the source. This problem can be viewed as the so-called CEO problem with additional secrecy constraints. We establish inner and outer bounds on the rate-distortion-equivocation region of this problem. We also obtain the region in special cases where the bounds are tight. Furthermore, we study the quadratic Gaussian case and provide the optimal rate-distortion-equivocation region when the eavesdropper has no side information and an achievable region for a more general setup with side information at the eavesdropper.

Index Terms—CEO problem, multiterminal source coding, secrecy constraints, eavesdropping, equivocation.

I. INTRODUCTION

As networks are becoming more distributed, their vulnerability to malicious activities increases which in turn raises the concern on the security of such networks. Consequently, information-theoretic security as a concrete framework for analyzing secrecy in networks has gained attention among researchers [2], [3]. Information-theoretic security, which was initially introduced by Shannon [4], exploits different statistical characteristics of received information at the legitimate receiver and at the eavesdropper. Moreover, it makes no assumptions on the computational power of the eavesdropper, unlike the traditional cryptographic approaches for secrecy. Later, Wyner introduced the Wiretap channel model in [5] and showed that perfectly secure communication without a shared secret key is possible if the channel from the transmitter to the eavesdropper is a degraded version of the channel to the legitimate receiver. This result was generalized to broadcast channels with confidential messages by Csiszár and Körner in [6]. Subsequently, many extensions to this problem have been developed and studied in the literature (see, for instance, [2], [3], and references therein).

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Authors are with the School of Electrical Engineering and ACCESS Linnaeus Center, KTH Royal Institute of Technology, SE-100 44 Stockholm, Sweden (emails: {naghibi,somayen,skoglund}@ee.kth.se).

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In this paper, we consider secrecy in a multiterminal source coding problem. In particular, we study the problem of conveying an information source to a single destination via multiple agents (encoders) in the presence of a passive eavesdropper. The agents have access to noisy observations of the source and are connected to the destination via noiseless rate-limited links. They do not cooperate or communicate to one another and are not required to estimate the source themselves. This scenario is of interest for many applications such as sensor networks or smart grid systems where reconstruction of the source at sensors and smart meters is not necessary. The distributed nature of such networks makes them more susceptible to eavesdropping. At each instant, the eavesdropper listens in on one of the links from the agents to the destination in order to obtain information about the source. In addition, it has access to side information correlated to the source. Since the link that will be compromised by the eavesdropper is unknown to the agents prior to their transmissions, each agent should protect its link in order to leak as little information as possible about the source. Our objective is to characterize the trade-off among agents' transmission rates, incurred distortion at the destination, and the amount of information revealed to the eavesdropper. This setup can be viewed as the extension of the so-called CEO problem [7] in which secrecy constraints are considered.

A. Related Work

The chief executive/estimation officer (CEO) problem was motivated in [7] by a communication and distributed processing system analogous to a scenario in which a firm's CEO is interested in information of a source that cannot be observed directly. The CEO assigns a group of agents to independently observe a corrupted version of the source and communicate their observations. The lossless variant of this setup was initially studied by Gel'fand and Pinsker [8]. It was extended by Yamamoto and Itoh [9] as well as Flynn and Gray [10] to the lossy case with only two encoders for which an achievable rate-distortion region was derived. The model was generalized to the CEO problem with many encoders by Berger and Viswanathan [7] in which the trade-off between the end-to-end average distortion and sum of the rates at which the agents transmit to the CEO was studied. Multiterminal lossy source coding problems, including the CEO problem, are still open in general. However, for the special case of the quadratic Gaussian CEO problem [11], the sum-rate-distortion function for infinite number of agents with identical signal-to-noise ratios (SNRs) was derived by Oohama [12], and later, the complete rate-distortion region with arbitrary number of agents and SNR values was characterized by Prabhakaran et al. [13].

and Oohama [14]. More recently, Courtade and Weissman [15] gave the rate-distortion region of the CEO problem under the logarithmic-loss distortion measure.

Secure lossless source coding with uncoded side information at the legitimate decoder and the eavesdropper was studied by Prabhakaran and Ramchandran [16] with the assumption of no rate constraint on the encoder-decoder link. The minimum leakage rate was derived and it was shown that due to the side information at the eavesdropper, the usual Slepian-Wolf scheme [17] is not always optimal. Lossless source coding with coded side information at the decoder (the so-called one-helper problem) and no side information at the eavesdropper was studied by Tandon et al. [18] where the rate-equivocation region was characterized. This setup was extended by Gündüz et al. [19] with additional side information at the eavesdropper in which inner and outer bounds on the compression-equivocation rate region were derived that did not match in general. Secure distributed lossless compression of two correlated sources, in which both sources were to be estimated at the decoder, was considered by Luh and Kundur [20] without side information at the eavesdropper and by Gündüz et al. [21] with side information at the eavesdropper. These models were generalized by Salimi et al. [22] to the case where both the legitimate receiver and the eavesdropper have access to correlated side information and the eavesdropper can choose to intercept either links from the encoders to the decoder at each instant. In [22], inner and outer bounds for the compression-equivocation region were provided which were proved to be tight for several special cases.

The extension to the lossy case was considered in [23]–[25], and more recently by Villard and Piantanida [26] in which inner and outer bounds on the rate-distortion-equivocation region were derived. The optimal characterization of the rate-distortion-equivocation region was first found in [24] for the lossy case with uncoded side information. Later in [26], the optimal characterization for the lossless case was also derived. A different setup was considered by Kittichokechai et al. [27] in which the eavesdropper can only access the coded side information, and the complete region was characterized under the logarithmic-loss distortion [15]. Chia and Kittichokechai [28] studied the case when the encoder has access to the side information of the decoder. Tandon et al. [29] considered a scenario with two legitimate receivers and investigated the privacy of side information at one receiver with respect to the other one. An alternative approach to provide secrecy in source coding problems is based on having a shared secret key between the transmitter and the legitimate receiver [30]–[32], although we do not exploit this approach in our work.

B. Contributions

Our setup in this paper has two main distinctions from the aforementioned scenarios; first, the destination (CEO) is interested in estimation of the original source rather than the agents' observations as in all prior works. Similarly, the secrecy constraints in our problem are on the equivocation of the eavesdropper with respect to the remote source, not to the observations of the agents. In fact, our setup is a generalization

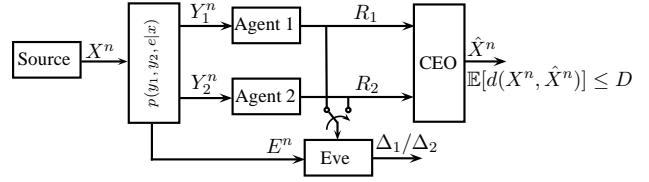


Fig. 1. The CEO problem with secrecy constraints.

of the previous cases considered for lossy secure source coding problems. We extend our previous work [33] for the lossless variant of this problem to the lossy case and derive inner and outer bounds on the rate-distortion-equivocation region of the CEO problem with secrecy constraints. We also investigate the region in special cases where the bounds are tight and we show that for these special cases our results coincide with the previous results in the literature.

In addition, we consider the quadratic Gaussian CEO problem with secrecy constraints and provide the optimal characterization of the rate-distortion-equivocation region for the case when the eavesdropper has no side information and an achievable region for a more general setup with side information at the eavesdropper.

C. Notations and Organization

In this paper, we use capital letters to indicate a random variable, small letters to indicate realization of a random variable, calligraphic letters to denote a set, e.g., \mathcal{X} , and $|\mathcal{X}|$ to indicate the cardinality of the set. The notation X^n denotes the sequence $\{X_1, \dots, X_n\}$. The notion $X - Y - Z$ shows that X , Y , and Z form a Markov chain, i.e., $p(x, y, z) = p(x, y)p(z|y)$ or $p(x, y, z) = p(x|y)p(y, z)$. We define $\mathcal{I}_M := \{1, \dots, M\}$ for $M \in \mathbb{N}$, and $[x]^+ := \max\{0, x\}$ for $x \in \mathbb{R}$. Finally, $\mathbb{1}_{\mathbb{R}_{>0}}(x) : \mathcal{X} \rightarrow \{0, 1\}$ denotes the indicator function such that $\mathbb{1}_{\mathbb{R}_{>0}}(x) = 1$ for $x \in \mathbb{R}_{>0}$, and $\mathbb{1}_{\mathbb{R}_{>0}}(x) = 0$ otherwise.

The rest of the paper is organized as follows: In Section II, we describe the problem along with some definitions. Main results for inner and outer bounds on the rate-distortion-equivocation region are presented in Section III. Then, we study some special cases of our results in Section IV where the region is completely characterized. The rate-distortion-equivocation region for the quadratic Gaussian case is given in Section V. Finally, the paper is concluded in Section VI.

II. PROBLEM SETTING

We consider the CEO problem with secrecy constraints as depicted in Fig. 1. In this setup, two non-cooperative and independent agents have access to length- n observations Y_1^n and Y_2^n , respectively, which are noisy versions of the source sequence X^n . These observations are conditionally independent given X^n . Each agent independently transmits a compressed version of its observation to the CEO over a rate-limited noiseless link. The CEO estimates the source sequence \hat{X}^n based on the received information from the two agents. An eavesdropper, referred to as Eve, with access to side information E^n correlated to the source sequence X^n can eavesdrop only one of the links from the agents to the CEO at

each time instance to obtain as much information as possible about the source. Therefore, agents' transmission rates should be such that the CEO can reconstruct the source reliably within a certain mean distortion threshold while simultaneously the equivocation at Eve is maximized. Eve's equivocation, with respect to either links, corresponds to her uncertainty about the original source when she combines her side information with the information obtained from the link. We assume that Eve cannot access both links simultaneously as the links are noise-free and in such case she would be more powerful than the CEO for estimating the source due to her additional side information. The sequences X^n, Y_1^n, Y_2^n , and E^n are independent and identically distributed (i.i.d.) according to joint distribution $p(x, y_1, y_2, e) = p(x)p(y_1|x)p(y_2|x)p(e|x)$ over the finite alphabet $\mathcal{X} \times \mathcal{Y}_1 \times \mathcal{Y}_2 \times \mathcal{E}$.

Let $d : \mathcal{X} \times \mathcal{X} \rightarrow [0, d_{\max}]$ be a finite distortion measure. We define the component-wise mean distortion between two sequences x^n, \hat{x}^n in \mathcal{X}^n as

$$d^{(n)}(x^n, \hat{x}^n) := \frac{1}{n} \sum_{i=1}^n d(x_i, \hat{x}_i). \quad (1)$$

Definition 1: A (M_1, M_2, n) -code for compression and transmission of the source by the agents with secrecy constraints consists of an encoding function at each agent, $f_j : \mathcal{Y}_j^n \rightarrow \mathcal{I}_{M_j}$ for $j \in \{1, 2\}$, and a decoding function at the CEO, $g : \mathcal{I}_{M_1} \times \mathcal{I}_{M_2} \rightarrow \mathcal{X}^n$. The equivocation rates for this code are defined as $\frac{1}{n} H(X^n | f_j(Y_j^n), E^n)$ for $j \in \{1, 2\}$.

Definition 2: A tuple $(R_1, R_2, \Delta_1, \Delta_2, D)$ is said to be achievable if $\forall \epsilon > 0$ there exists $N_0 \in \mathbb{N}_{>0}$ such that for all $n > N_0$ there exists a sequence of (M_1, M_2, n) -codes with

$$\begin{aligned} \log(M_j) &\leq n(R_j + \epsilon), \quad \forall j \in \{1, 2\}, \\ H(X^n | f_j(Y_j^n), E^n) &\geq n(\Delta_j - \epsilon), \quad \forall j \in \{1, 2\}, \\ \mathbb{E}[d^{(n)}(X^n, g(f_1(Y_1^n), f_2(Y_2^n)))] &\leq D + \epsilon. \end{aligned}$$

Let \mathcal{R} denote the rate-distortion-equivocation region defined as the set of all achievable tuples $(R_1, R_2, \Delta_1, \Delta_2, D)$.

III. INNER AND OUTER BOUNDS ON THE RATE-DISTORTION-EQUIVOCATION REGION

A. Inner Bound

Theorem 1: Let \mathcal{R}_{in} denote the region defined as the closure of the convex hull of the set of all tuples $(R_1, R_2, \Delta_1, \Delta_2, D)$ such that there exist random variables V_1, V_2, U_1 , and U_2 on some finite sets $\mathcal{V}_1, \mathcal{V}_2, \mathcal{U}_1$, and \mathcal{U}_2 , respectively, according to the distribution $p(x, y_1, y_2, e, v_1, v_2, u_1, u_2) = p(x)p(y_1|x)p(y_2|x)p(e|x)p(u_1|y_1)p(u_2|y_2)p(v_1|u_1)p(v_2|u_2)$ and a function $\hat{X} : \mathcal{U}_1 \times \mathcal{U}_2 \rightarrow \mathcal{X}$ that satisfy

$$R_1 \geq I(U_1; Y_1|U_2), \quad (2)$$

$$R_2 \geq I(U_2; Y_2|U_1), \quad (3)$$

$$R_1 + R_2 \geq I(U_1, U_2; Y_1, Y_2), \quad (4)$$

$$\begin{aligned} \Delta_1 &\leq [H(X|V_1, E) - I(U_1; Y_1|V_1, U_2) \\ &\quad + I(U_1; Y_1|V_1, X)]^+, \end{aligned} \quad (5)$$

$$\begin{aligned} \Delta_2 &\leq [H(X|V_2, E) - I(U_2; Y_2|V_2, U_1) \\ &\quad + I(U_2; Y_2|V_2, X)]^+, \end{aligned} \quad (6)$$

$$\begin{aligned} \Delta_1 + \Delta_2 &\leq [H(X|V_1, E) + H(X|V_2, E) \\ &\quad - I(U_1, U_2; Y_1, Y_2|V_1, V_2) \\ &\quad + I(U_1; Y_1|V_1, X) + I(U_2; Y_2|V_2, X)]^+, \end{aligned} \quad (7)$$

$$\begin{aligned} \Delta_1 - R_2 &\leq [H(X|V_1, E) - I(U_2; Y_2|U_1) \\ &\quad - I(U_1; Y_1|V_1) + I(U_1; Y_1|V_1, X)]^+, \end{aligned} \quad (8)$$

$$\begin{aligned} \Delta_2 - R_1 &\leq [H(X|V_2, E) - I(U_1; Y_1|U_2) \\ &\quad - I(U_2; Y_2|V_2) + I(U_2; Y_2|V_2, X)]^+, \end{aligned} \quad (9)$$

$$D \geq \mathbb{E}[d(X, \hat{X}(U_1, U_2))]. \quad (10)$$

Then, we have $\mathcal{R}_{\text{in}} \subset \mathcal{R}$.

Proof: The proof is given in Appendix A. \blacksquare

Proposition 1: In Theorem 1, it suffices to consider auxiliary random variables V_j and U_j for $j \in \{1, 2\}$ with cardinalities $|\mathcal{V}_j| \leq |\mathcal{Y}_j| + 9$ and $|\mathcal{U}_j| \leq (|\mathcal{Y}_j| + 9)(|\mathcal{Y}_j| + 5)$, respectively (see Appendix B for the proof).

The achievability scheme resulting in the inner bound is based on superposition coding and random binning at the agents, and joint decoding at the CEO. In particular, agent j first transmits the bin index related to the auxiliary random variable V_j with distribution $p(v_j|u_j)$ via the noiseless link. Then, the agents send the remaining information which is required for the CEO to be able to reconstruct the source based on the Wyner-Ziv scheme [34]. The detailed proof is given in Appendix A, however, we provide some intuitions on the results. Inequalities (2)–(4) and (10) are similar to the Berger-Tung bounds [35], [36] that establish perfect estimation of U_1 and U_2 at the CEO from which X can be reconstructed within the distortion limit D . In the equivocation bounds (5) and (6), the first term corresponds to Eve's uncertainty about the source after decoding the codeword v_j^n based on the received bin index combined with her side information and the second term is the reduction in her uncertainty when receiving the remaining information transmitted to the CEO by the agents. Finally, the last term in (5) and (6) stems from the fact that in contrast to previous works, the secrecy constraints are on Eve's equivocation with respect to the original source while the transmitted information by the agents are functions of their respective observations and not the source, resulting in an increase in Eve's uncertainty. Inequalities (8) and (9) depict a trade-off between Eve's equivocation and transmission rates, implying that each link's transmission rate limits the other link's equivocation rate.

Remark 1: The region of Theorem 1 can also be obtained by constructing six different codes achieving the corner points shown in Tables I–II and using the time-sharing technique between these points. Each corner point is achieved using a four-step communication to transmit variables V_1, V_2, U_1 , and U_2 to the CEO with different decoding orders, provided that V_j is decoded prior to U_j for $j \in \{1, 2\}$. In each step, previously received and decoded information at the CEO is used as side information for the current decoding step. Each code employs superposition coding, with V_j as the first layer and U_j as the second layer, and random binning based on the available side information at the CEO in each communication step.

TABLE I
CORNER POINTS OF THE INNER REGION CORRESPONDING TO DIFFERENT DECODING ORDERS: RATES AND DISTORTION.

Corner point	Decoding order	R_1	R_2	D
1	V_2, U_2, V_1, U_1	$I(U_1; Y_1 U_2)$	$I(U_2; Y_2)$	$\mathbb{E}[d(X, \hat{X}(U_1, U_2))]$
2	V_2, V_1, U_2, U_1	$I(V_1; Y_1 V_2) + I(U_1; Y_1 V_1, U_2)$	$I(V_2; Y_2) + I(U_2; Y_2 V_1, V_2)$	$\mathbb{E}[d(X, \hat{X}(U_1, U_2))]$
3	V_1, V_2, U_2, U_1	$I(V_1; Y_1) + I(U_1; Y_1 V_1, U_2)$	$I(U_2; Y_2 V_1)$	$\mathbb{E}[d(X, \hat{X}(U_1, U_2))]$
4	V_1, U_1, V_2, U_2	$I(U_1; Y_1)$	$I(U_2; Y_2 U_1)$	$\mathbb{E}[d(X, \hat{X}(U_1, U_2))]$
5	V_1, V_2, U_1, U_2	$I(V_1; Y_1) + I(U_1; Y_1 V_1, V_2)$	$I(V_2; Y_2 V_1) + I(U_2; Y_2 V_2, U_1)$	$\mathbb{E}[d(X, \hat{X}(U_1, U_2))]$
6	V_2, V_1, U_1, U_2	$I(U_1; Y_1 V_2)$	$I(V_2; Y_2) + I(U_2; Y_2 V_2, U_1)$	$\mathbb{E}[d(X, \hat{X}(U_1, U_2))]$

TABLE II
CORNER POINTS OF THE INNER REGION CORRESPONDING TO DIFFERENT DECODING ORDERS: EQUIVOCATION RATES.

Corner point	Decoding order	Δ_1	Δ_2
1	V_2, U_2, V_1, U_1	$[H(X V_1, E) - I(U_1; Y_1 V_1, U_2) + I(U_1; Y_1 V_1, X)]^+$	$[H(X V_2, E) - I(U_2; Y_2 V_2) + I(U_2; Y_2 V_2, X)]^+$
2	V_2, V_1, U_2, U_1	$[H(X V_1, E) - I(U_1; Y_1 V_1, U_2) + I(U_1; Y_1 V_1, X)]^+$	$[H(X V_2, E) - I(U_2; Y_2 V_1, V_2) + I(U_2; Y_2 V_2, X)]^+$
3	V_1, V_2, U_2, U_1	$[H(X V_1, E) - I(U_1; Y_1 V_1, U_2) + I(U_1; Y_1 V_1, X)]^+$	$[H(X V_2, E) - I(U_2; Y_2 V_1, V_2) + I(U_2; Y_2 V_2, X)]^+$
4	V_1, U_1, V_2, U_2	$[H(X V_1, E) - I(U_1; Y_1 V_1) + I(U_1; Y_1 V_1, X)]^+$	$[H(X V_2, E) - I(U_2; Y_2 V_2, U_1) + I(U_2; Y_2 V_2, X)]^+$
5	V_1, V_2, U_1, U_2	$[H(X V_1, E) - I(U_1; Y_1 V_1, V_2) + I(U_1; Y_1 V_1, X)]^+$	$[H(X V_2, E) - I(U_2; Y_2 V_2, U_1) + I(U_2; Y_2 V_2, X)]^+$
6	V_2, V_1, U_1, U_2	$[H(X V_1, E) - I(U_1; Y_1 V_1, V_2) + I(U_1; Y_1 V_1, X)]^+$	$[H(X V_2, E) - I(U_2; Y_2 V_2, U_1) + I(U_2; Y_2 V_2, X)]^+$

B. Outer Bound

Theorem 2: Let \mathcal{R}_{out} denote the region defined as the closure of the set of all tuples $(R_1, R_2, \Delta_1, \Delta_2, D)$ such that there exist random variables V_1, V_2, U_1 , and U_2 on some finite sets $\mathcal{V}_1, \mathcal{V}_2, \mathcal{U}_1$, and \mathcal{U}_2 , respectively, which form Markov chains $V_j - U_j - Y_j - (X, E, Y_{j'})$ for $j, j' \in \{1, 2\}$ with $j \neq j'$, and a function $\hat{X} : \mathcal{U}_1 \times \mathcal{U}_2 \rightarrow \mathcal{X}$ that satisfy

$$R_1 \geq I(U_1; Y_1|U_2), \quad (11)$$

$$R_2 \geq I(U_2; Y_2|U_1), \quad (12)$$

$$R_1 + R_2 \geq I(U_1, U_2; Y_1, Y_2), \quad (13)$$

$$\Delta_1 \leq H(X|E) - I(X; V_1|E), \quad (14)$$

$$\Delta_2 \leq H(X|E) - I(X; V_2|E), \quad (15)$$

$$\Delta_1 - R_2 \leq H(X|E) - I(X; V_1|E) - I(X; V_2|V_1, E), \quad (16)$$

$$\Delta_2 - R_1 \leq H(X|E) - I(X; V_2|E) - I(X; V_1|V_2, E), \quad (17)$$

$$D \geq \mathbb{E}[d(X, \hat{X}(U_1, U_2))]. \quad (18)$$

Then, we have $\mathcal{R} \subset \mathcal{R}_{\text{out}}$.

Proof: The proof is given in Appendix C. \blacksquare

Proposition 2: In Theorem 2, it suffices to consider auxiliary random variables V_j and U_j for $j \in \{1, 2\}$ with cardinalities $|\mathcal{V}_j| \leq |\mathcal{Y}_j| + 7$ and $|\mathcal{U}_j| \leq (|\mathcal{Y}_j| + 7)(|\mathcal{Y}_j| + 3)$, respectively (see Appendix D for the proof).

IV. SPECIAL CASE: THE ONE-HELPER PROBLEM WITH SECRECY CONSTRAINTS

If Agent 1 has access to the source sequence X^n , our setup reduces to the lossy source coding problem with a helper and an eavesdropper who can choose to listen in on either source-destination or helper-destination links.

Corollary 1: In the above setting, if we additionally assume the helper's link is perfectly secure, our results coincide with the results given by Villard and Piantanida [26, Theorem 3].

The inner bound is obtained by setting $Y_1 = X$, $V_2 = U_2$, and removing the constraints on Δ_2 in Theorem 1, and the outer bound can be proved similar to the proof of Theorem 2.

Corollary 2: In the described one-helper problem with secrecy constraints, if $R_2 \geq H(Y_2)$, the helper's sequence can be reconstructed by the destination losslessly. Then, the rate-distortion-equivocation region is characterized by

$$R_1 \geq I(X; U_1|Y_2), \quad (19)$$

$$\Delta_1 \leq [H(X|V_1, E) - I(U_1; Y_1|V_1, Y_2)]^+, \quad (20)$$

$$\Delta_2 \leq H(X|Y_2, E), \quad (21)$$

$$D \geq \mathbb{E}[d(X, \hat{X}(U_1, Y_2))], \quad (22)$$

where the auxiliary random variables V_1 and U_1 satisfy the Markov chain $V_1 - U_1 - X - (E, Y_2)$.

The achievability proof follows from the proof of Theorem 1 by setting $Y_1 = X$ and $V_2 = U_2 = Y_2$. Inequalities (7)–(9) are inactive for this setup. The converse proof is given in [26, Theorem 3] for the secure lossy source coding with uncoded side information. Note that if Eve intercepts the helper's link, it can also reconstruct the helper's sequence Y_2^n losslessly.

Corollary 3: For the lossless one-helper setting, i.e., $D = 0$, if $R_2 \geq H(Y_2)$, the rate-equivocation region writes as:

$$R_1 \geq H(X|Y_2), \quad (23)$$

$$\Delta_1 \leq I(X; Y_2|V_1) - I(X; E|V_1), \quad (24)$$

$$\Delta_2 \leq H(X|Y_2E). \quad (25)$$

The achievability proof follows from the proof of Theorem 1 by setting $U_1 = Y_1 = X$ and $V_2 = U_2 = Y_2$. The converse proof is similar to the proof given in [16].

Corollary 4: For the lossless one-helper setting, i.e., $D = 0$, if the eavesdropper has no side information, the rate-

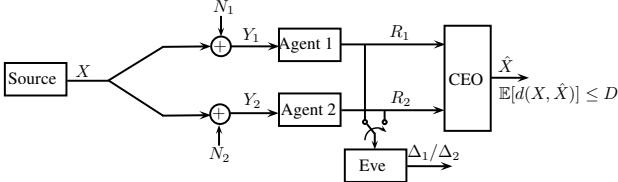


Fig. 2. The quadratic Gaussian case with no side information at Eve.

equivocation region is characterized by

$$R_1 \geq H(X|U_2), \quad (26)$$

$$R_2 \geq I(Y_2; U_2), \quad (27)$$

$$\Delta_1 \leq I(X; U_2), \quad (28)$$

$$\Delta_2 \leq H(X|U_2). \quad (29)$$

The achievability is a special case of Theorem 1 and obtained by setting V_1 and E to be constants, $U_1 = Y_1 = X$, and $V_2 = U_2$. The proof of converse is given in Appendix E.

V. THE QUADRATIC GAUSSIAN CASE

In this section, we study the Gaussian CEO problem with secrecy constraints and quadratic distortion measure.

Let X be a Gaussian source, i.e., $X \sim \mathcal{N}(0, \sigma_X^2)$. The observations at the agents are modeled as $Y_j = X + N_j$ for $j \in \{1, 2\}$, with $N_j \sim \mathcal{N}(0, \sigma_{N_j}^2)$, where Gaussian random variables X , N_1 , and N_2 are mutually independent.

First, we consider the case where the eavesdropper has no side information. The model is depicted in Fig. 2 and the following theorem provides the complete rate-distortion-equivocation region for this Gaussian setup.

Theorem 3: In the quadratic Gaussian CEO problem with secrecy constraints, the rate-distortion-equivocation region is characterized by the set of all tuples $(R_1, R_2, \Delta_1, \Delta_2, D)$ satisfying

$$R_1 \geq r_1 + \frac{1}{2} \log \frac{1}{D} - \frac{1}{2} \log \left(\frac{1}{\sigma_X^2} + \frac{1 - 2^{-2r_2}}{\sigma_{N_2}^2} \right), \quad (30)$$

$$R_2 \geq r_2 + \frac{1}{2} \log \frac{1}{D} - \frac{1}{2} \log \left(\frac{1}{\sigma_X^2} + \frac{1 - 2^{-2r_1}}{\sigma_{N_1}^2} \right), \quad (31)$$

$$R_1 + R_2 \geq r_1 + r_2 + \frac{1}{2} \log \frac{1}{D} - \frac{1}{2} \log \frac{1}{\sigma_X^2}, \quad (32)$$

$$\begin{aligned} \Delta_1 &\leq \frac{1}{2} \log(2\pi e \sigma_X^2) - \frac{1}{2} \log \frac{1}{D} \\ &\quad + \frac{1}{2} \log \left(\frac{1}{\sigma_X^2} + \frac{1 - 2^{-2r_2}}{\sigma_{N_2}^2} \right), \end{aligned} \quad (33)$$

$$\begin{aligned} \Delta_2 &\leq \frac{1}{2} \log(2\pi e \sigma_X^2) - \frac{1}{2} \log \frac{1}{D} \\ &\quad + \frac{1}{2} \log \left(\frac{1}{\sigma_X^2} + \frac{1 - 2^{-2r_1}}{\sigma_{N_1}^2} \right), \end{aligned} \quad (34)$$

$$\Delta_1 + \Delta_2 \leq \log(2\pi e) - \frac{1}{2} \log \frac{1}{D} - \frac{1}{2} \log \frac{1}{\sigma_X^2}, \quad (35)$$

$$\Delta_1 - R_2 \leq \frac{1}{2} \log(2\pi e) - \frac{1}{2} \log \frac{1}{D} - r_2, \quad (36)$$

$$\Delta_2 - R_1 \leq \frac{1}{2} \log(2\pi e) - \frac{1}{2} \log \frac{1}{D} - r_1, \quad (37)$$

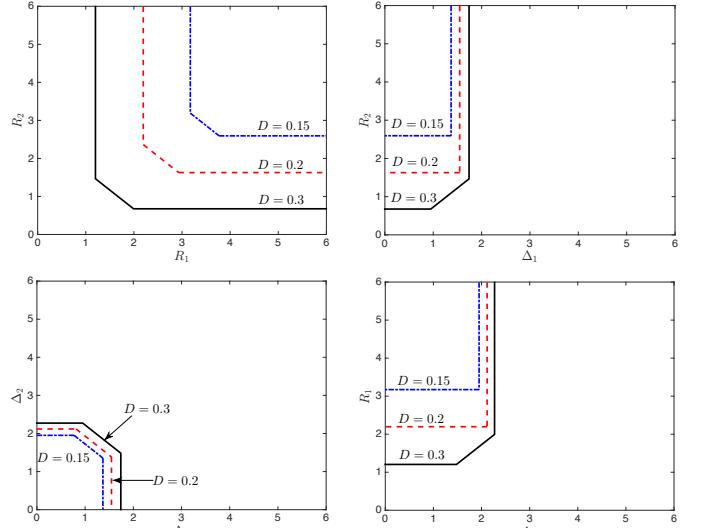


Fig. 3. An example of the rate-distortion-equivocation region for the quadratic Gaussian case with no side information at Eve and different distortion constraints.

for some $(r_1, r_2) \in \mathbb{R}_{\geq 0}^2$ that satisfy

$$\frac{1}{D} \leq \frac{1}{\sigma_X^2} + \frac{1 - 2^{-2r_1}}{\sigma_{N_1}^2} + \frac{1 - 2^{-2r_2}}{\sigma_{N_2}^2}. \quad (38)$$

Proof: The proof is given in Appendix F. ■

An example of the region of Theorem 3 is illustrated in Fig. 3 for different distortion constraints.

Remark 2: We note that equivocation as the secrecy measure in the finite alphabet setting represents the uncertainty of Eve about the source, but in the Gaussian setting, this interpretation is not quite valid. However, based on [37, Theorem 8.6.6], we can relate the equivocation rate (normalized differential entropy) to the estimation error at the eavesdropper. That is, we define the secrecy measure in the Gaussian setting as

$$\frac{1}{n} \mathbb{E} \left[\|X^n - \hat{Z}^{(n)}(f_j(Y_j^n), E^n)\|^2 \right] \geq \frac{1}{2\pi e} e^{2\Delta_j} \triangleq \theta_j, \quad (39)$$

for $j = 1, 2$, where $\hat{Z}^{(n)}(f_j(Y_j^n), E^n)$ is an estimator of X^n at the eavesdropper. Then, we have $\Delta_j = \frac{1}{2} \log(2\pi e \theta_j)$, and in this sense, equivocation rate provides a lower bound on the normalized distortion at Eve. Moreover, we can relate this to the information leakage rate as

$$\begin{aligned} \frac{1}{n} I(X^n; f_j(Y_j^n, E^n)) &= h(X) - \frac{1}{n} h(f_j(Y_j^n, E^n)) \\ &\leq h(X) - \Delta_j \\ &= h(X) - \frac{1}{2} \log(2\pi e \theta_j), \end{aligned} \quad (40)$$

which is also in line with the result in [38, Theorem 7.3].

Next, we consider the case where Eve has access to additional side information correlated to the source as shown in Fig. 4. We model this side information as $E = X + N_E$ where N_E is a Gaussian random variable with $N_E \sim \mathcal{N}(0, \sigma_{N_E}^2)$ and is independent of X , N_1 , and N_2 . The following theorem gives an inner bound for the rate-distortion-equivocation region of the quadratic Gaussian CEO problem with secrecy constraints and side information at the eavesdropper.

Theorem 4: In the quadratic Gaussian CEO problem with secrecy constraints and side information at the eavesdropper, a tuple $(R_1, R_2, \Delta_1, \Delta_2, D)$ is achievable if

$$R_1 \geq r_1 + \frac{1}{2} \log \frac{1}{D} - \frac{1}{2} \log \left(\frac{1}{\sigma_X^2} + \frac{1 - 2^{-2r_2}}{\sigma_{N_2}^2} \right), \quad (41)$$

$$R_2 \geq r_2 + \frac{1}{2} \log \frac{1}{D} - \frac{1}{2} \log \left(\frac{1}{\sigma_X^2} + \frac{1 - 2^{-2r_1}}{\sigma_{N_1}^2} \right), \quad (42)$$

$$R_1 + R_2 \geq r_1 + r_2 + \frac{1}{2} \log \frac{1}{D} - \frac{1}{2} \log \frac{1}{\sigma_X^2}, \quad (43)$$

$$\begin{aligned} \Delta_1 &\leq \frac{1}{2} \log(2\pi e) - \frac{1}{2} \log \left(\frac{1}{\sigma_X^2} + \frac{1}{\sigma_{N_E}^2} \right) \\ &\quad - \frac{1}{2} \log \frac{1}{D} + \frac{1}{2} \log \left(\frac{1}{\sigma_X^2} + \frac{1 - 2^{-2r_2}}{\sigma_{N_E}^2} \right) + T_1, \quad (44) \end{aligned}$$

$$\Delta_2 \leq \frac{1}{2} \log(2\pi e) - \frac{1}{2} \log \left(\frac{1}{\sigma_X^2} + \frac{1}{\sigma_{N_E}^2} \right) - \frac{1}{2} \log \frac{1}{D} + \frac{1}{2} \log \left(\frac{1}{\sigma_X^2} + \frac{1 - 2^{-2r_1}}{\sigma_{N_1}^2} \right) + T_2, \quad (45)$$

$$\begin{aligned}
\Delta_1 + \Delta_2 &\leq \log(2\pi e) - \log\left(\frac{1}{\sigma_X^2} + \frac{1}{\sigma_{N_E}^2}\right) - \frac{1}{2} \log \frac{1}{D} \\
&\quad + \frac{1}{2} \log \frac{1}{\sigma_X^2} + T_1 + T_2 \\
&\quad + \frac{1}{2} \log \left(\frac{\frac{\sigma_X^2 + \frac{\sigma_{N_1}^2}{1-2^{-2r_1}}}{\sigma_{N_1}^2}}{\frac{1-2^{-2r_1}}{1-2^{-2r_2}} + \frac{\sigma_{N_2}^2}{1-2^{-2r_2}}} \right) \mathbb{1}_{\mathbb{R}_{>0}}(T_1 + T_2),
\end{aligned} \tag{46}$$

$$\begin{aligned} \Delta_1 - R_2 &\leq \frac{1}{2} \log(2\pi e) - \frac{1}{2} \log \left(\frac{1}{\sigma_X^2} + \frac{1}{\sigma_{N_E}^2} \right) \\ &\quad - \frac{1}{2} \log \frac{1}{D} - r_2 + T_1 \\ &\quad + \frac{1}{2} \log \left(\frac{\sigma_X^2 + \frac{\sigma_{N_1}^2}{1-2^{-2r_1}}}{\frac{\sigma_{N_1}^2}{1-2^{-2r_1}} + \frac{\sigma_{N_2}^2}{1-2^{-2r_2}}} \right) \mathbb{1}_{\mathbb{R}_{>0}}(T_1), \end{aligned} \quad (47)$$

$$\begin{aligned} \Delta_2 - R_1 &\leq \frac{1}{2} \log(2\pi e) - \frac{1}{2} \log \left(\frac{1}{\sigma_X^2} + \frac{1}{\sigma_{N_E}^2} \right) \\ &\quad - \frac{1}{2} \log \frac{1}{D} - r_1 + T_2 \\ &\quad + \frac{1}{2} \log \left(\frac{\sigma_X^2 + \frac{\sigma_{N_1}^2}{1-2^{-2r_1}}}{\frac{\sigma_{N_1}^2}{1-2^{-2r_1}} + \frac{\sigma_{N_2}^2}{1-2^{-2r_2}}} \right) \mathbb{1}_{\mathbb{R}_{>0}}(T_2), \end{aligned} \quad (48)$$

where $1_{\mathbb{R}_{>0}}(\cdot)$ is the indicator function and

$$T_1 := \max \left\{ 0, \frac{1}{2} \log \left(1 + \frac{\frac{\sigma_{N_2}^2}{1-2^{-2T_2}} - \sigma_{N_E}^2}{\frac{\sigma_{N_1}^2}{1-2^{-2T_1}} + \sigma_{N_{-}}^2} \right) \right\}, \quad (49)$$

$$T_2 := \max \left\{ 0, \frac{1}{2} \log \left(1 + \frac{\frac{\sigma_{N_1}^2}{1-2^{-2}r_1} - \sigma_{N_E}^2}{\frac{\sigma_{N_2}^2}{1-2^{-2}r_2} + \sigma_{N_E}^2} \right) \right\}, \quad (50)$$

for some $(r_1, r_2) \in \mathbb{R}_{\geq 0}^2$ that satisfy

$$\frac{1}{D} \leq \frac{1}{\sigma_{\Sigma}^2} + \frac{1 - 2^{-2r_1}}{\sigma_{\Sigma}^2} + \frac{1 - 2^{-2r_2}}{\sigma_{\Sigma}^2}. \quad (51)$$

Proof: The proof is given in Appendix G.

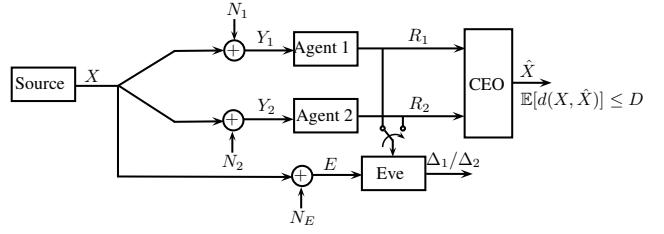


Fig. 4. The quadratic Gaussian case with side information at Eve.

Note that if there is no correlation between Eve's side information and the source, i.e., $\sigma_{N_E}^2 \rightarrow \infty$, the region of Theorem 4 coincides with the one in Theorem 3.

VI. CONCLUSION

We studied the extension of the CEO problem with secrecy constraints. This setup is of interest to communication scenarios such as sensor networks or smart power grids in which links are vulnerable to eavesdropping. We derived inner and outer bounds on the rate-distortion-equivocation region in the discrete case. We also showed that the results that were derived for the one-helper problem with secrecy constraints in [18] and [26] can be obtained as special cases of our results for the CEO problem with secrecy constraints. In addition, we provide the optimal region for the quadratic Gaussian case when Eve has no side information as well as an achievable region for a more general case. In this work, we have considered noise-free links from the agents to the CEO, however, it would be interesting to investigate the effects of noisy channels in this problem. Moreover, extending this problem to include more agents and eavesdroppers with possibly different side information is another direction worthwhile investigating.

APPENDIX A

PROOF OF THEOREM 1: THE INNER BOUND

We first state the following lemma that we use in the proof of Theorem 1. The lemma follows from [2, Section 2.3].

Lemma 1: Let \mathcal{C}_n be a random codebook and s be the corresponding codeword index of V^n from this codebook. Let $\Pr((V^n(s), E^n) \in \mathcal{T}_\epsilon^{(n)}) \rightarrow 1$ as $n \rightarrow \infty$, where $\mathcal{T}_\epsilon^{(n)}$ denotes the set of jointly ϵ -typical n -sequences. Then, we have

$$H(E^n|s, \mathcal{C}_n) \leq n(H(E|V) + \epsilon). \quad (52)$$

Proof of Lemma 1: Let Z be a binary variable such that

$$Z = \begin{cases} 1 & \text{if } (V^n, E^n) \in \mathcal{T}_\epsilon^{(n)} \\ 0 & \text{otherwise.} \end{cases}$$

If $\Pr(Z = 0) \leq \epsilon$ for sufficiently large n , we have

$$\begin{aligned}
H(E^n|s, \mathcal{C}_n) &\leq H(E^n|V^n, \mathcal{C}_n) \\
&\leq H(E^n|V^n, Z) + H(Z) \\
&\leq \Pr(Z=1)H(E^n|V^n, Z=1) \\
&\quad + \Pr(Z=0)H(E^n|V^n, Z=0) + H(Z) \\
&\leq H(E^n|V^n, Z=1) + n\epsilon \log |\mathcal{E}| + H(Z) \\
&\leq \sum_{v^n \in \mathcal{I}_\epsilon^{(n)}} p(v^n|Z=1)H(E^n|V^n = v^n, Z=1)
\end{aligned}$$

$$\begin{aligned}
& + n\epsilon \log |\mathcal{E}| + H(Z) \\
\leq & \sum_{v^n \in \mathcal{T}_\epsilon^{(n)}} p(v^n|Z=1) \log |\mathcal{T}_\epsilon^{(n)}(E|v^n)| \\
& + n\epsilon \log |\mathcal{E}| + H_2(\epsilon) \\
\leq & n(H(E|V) + \epsilon'), \tag{53}
\end{aligned}$$

where $H_2(\cdot)$ is the binary entropy function. The last inequality is due to properties of jointly typical sequences [39, Chapter 2]. \blacksquare

Now, we proceed to prove Theorem 1.

Let V_1 , V_2 , U_1 , and U_2 be random variables on some finite sets \mathcal{V}_1 , \mathcal{V}_2 , \mathcal{U}_1 , and \mathcal{U}_2 according to the joint distribution $p(x, y_1, y_2, e, v_1, v_2, u_1, u_2) = p(x)p(y_1|x)p(y_2|x)p(e|x)p(u_1|y_1)p(u_2|y_2)p(v_1|u_1)p(v_2|u_2)$, along with a function $\hat{X} : \mathcal{U}_1 \times \mathcal{U}_2 \rightarrow \mathcal{X}$ satisfying the conditions of Theorem 1.

Codebook generation: For fixed conditional distributions $p(u_j|y_j)$ and $p(v_j|u_j)$, $j = 1, 2$, randomly generate $2^{n(I(V_j;Y_j)+\epsilon_1)}$ independent codewords $v_j^n(s_j)$ of length n according to $\prod_{i=1}^n P_{V_j}(v_{j,i}(s_j))$, where $s_j \in \{1, \dots, 2^{n(I(V_j;Y_j)+\epsilon_1)}\}$. Then, divide them into $2^{n(R_{V_j}+\epsilon_2)}$ equal-sized bins, indexed by $b_j \in \{1, \dots, 2^{n(R_{V_j}+\epsilon_2)}\}$ and denoted by $\{\mathcal{B}_j(b_j)\}$. For each codeword $v_j^n(s_j)$, randomly generate $2^{n(I(U_j;Y_j)+\epsilon_3)}$ independent sequences $u_j^n(s_j, s'_j)$ according to $\prod_{i=1}^n P_{U_j}(u_{j,i}(s_j, s'_j))$, and divide them into $2^{n(R_{U_j}+\epsilon_4)}$ equal-sized bins, indexed by $w_j \in \{1, \dots, 2^{n(R_{U_j}+\epsilon_4)}\}$ and denoted by $\{\mathcal{B}'_j(s_j, w_j)\}$. Define $R_j = R_{V_j} + R_{U_j}$ for $j = 1, 2$. The codebook is revealed to the agents, CEO, and Eve.

Encoding: Assume that the sequence y_j^n is observed by Agent j , $j = 1, 2$. Find a codeword $v_j^n(s_j)$ jointly typical with y_j^n . If there is more than one such codeword, select one uniformly at random. If there is no such v_j^n , select one out of $2^{n(I(V_j;Y_j)+\epsilon_1)}$ uniformly at random. Given $v_j^n(s_j)$, find a codeword $u_j^n(s_j, s'_j)$ jointly typical with y_j^n . If there is more than one such codeword, select one uniformly at random. If there is no such u_j^n , select one out of $2^{n(I(U_j;Y_j)+\epsilon_3)}$ uniformly at random. The agent transmits the bin indices b_j and w_j of the codewords $v_j^n(s_j) \in \mathcal{B}_j(b_j)$ and $u_j^n(s_j, s'_j) \in \mathcal{B}'_j(s_j, w_j)$, respectively, i.e., $f_j(y_j^n) = (b_j, w_j)$.

Decoding at the CEO: Given the received messages from both agents, $J_1 = (b_1, w_1)$ and $J_2 = (b_2, w_2)$, find a unique index tuple $(\hat{s}_1, \hat{s}'_1, \hat{s}_2, \hat{s}'_2)$ such that the codewords $(v_1^n(\hat{s}_1), u_1^n(\hat{s}_1, \hat{s}'_1), v_2^n(\hat{s}_2), u_2^n(\hat{s}_2, \hat{s}'_2))$ are jointly typical, and they are in the bin indexed by (b_1, w_1, b_2, w_2) . If there is such a unique index tuple, compute the source estimate component-wise as $\hat{x}_i = g_i(J_1, J_2) := \hat{X}(u_{1,i}(\hat{s}_1, \hat{s}'_1), u_{2,i}(\hat{s}_2, \hat{s}'_2))$ for $i = 1, \dots, n$; otherwise set the output to an arbitrary sequence in \mathcal{X}^n .

Error analysis: Let (s_1, s'_1, s_2, s'_2) and $(\hat{s}_1, \hat{s}'_1, \hat{s}_2, \hat{s}'_2)$ be the chosen indices at the encoders and decoder, respectively. Let $\Pr(\mathbf{E})$ denote the probability of an error event during encoding and decoding steps. We now show that this probability, averaged over all possible codebooks, tends to zero as $n \rightarrow \infty$ provided that conditions of Theorem 1 is satisfied. Consider the following error events in the encoding steps (for

$j, j' = 1, 2$, and $j \neq j'$):

$$\begin{aligned}
\mathbf{E}_0 &= \left\{ (S^n, Y_1^n, Y_2^n, E^n) \notin \mathcal{T}_\delta^{(n)}(S, Y_1, Y_2, E) \right\}, \\
\mathbf{E}_1 &= \mathbf{E}_0^c \cap \left\{ (V_j^n(\tilde{s}_j), Y_j^n) \notin \mathcal{T}_\delta^{(n)}(V_j, Y_j) \right\}, \\
&\quad \forall \tilde{s}_j \in \{1, \dots, 2^{n(I(V_j;Y_j)+\epsilon_1)}\}, \\
\mathbf{E}_2 &= \mathbf{E}_0^c \cap \mathbf{E}_1^c \cap \left\{ (U_j^n(s_j, \tilde{s}'_j), Y_j^n) \notin \mathcal{T}_\delta^{(n)}(U_j, Y_j | v_j^n(s_j)) \right\}, \\
&\quad \forall \tilde{s}'_j \in \{1, \dots, 2^{n(I(U_j;Y_j)+\epsilon_3)}\}, \\
\mathbf{E}_3 &= \left(\bigcap_{t=0}^2 \mathbf{E}_t^c \right) \cap \left\{ (V_j^n(s_j), Y_1^n, Y_2^n) \notin \mathcal{T}_\delta^{(n)}(V_j, Y_1, Y_2) \right\}, \\
\mathbf{E}_4 &= \left(\bigcap_{t=0}^3 \mathbf{E}_t^c \right) \cap \left\{ (U_j^n(s_j, s'_j), Y_1^n, Y_2^n) \right. \\
&\quad \left. \notin \mathcal{T}_\delta^{(n)}(U_j, Y_1, Y_2 | v_j^n(s_j)) \right\}, \\
\mathbf{E}_5 &= \left(\bigcap_{t=0}^4 \mathbf{E}_t^c \right) \cap \left\{ (V_j^n(s_j), U_{j'}^n(s_{j'}, s'_{j'}), Y_1^n, Y_2^n) \right. \\
&\quad \left. \notin \mathcal{T}_\delta^{(n)}(V_j, U_{j'}, Y_1, Y_2 | v_{j'}^n(s_{j'})) \right\}, \\
\mathbf{E}_6 &= \left(\bigcap_{t=0}^5 \mathbf{E}_t^c \right) \cap \left\{ (V_1^n(s_1), V_2^n(s_2), Y_1^n, Y_2^n) \right. \\
&\quad \left. \notin \mathcal{T}_\delta^{(n)}(V_1, V_2, Y_1, Y_2) \right\}, \\
\mathbf{E}_7 &= \left(\bigcap_{t=0}^6 \mathbf{E}_t^c \right) \cap \left\{ (U_1^n(s_1, s'_1), U_2^n(s_2, s'_2), Y_1^n, Y_2^n) \right. \\
&\quad \left. \notin \mathcal{T}_\delta^{(n)}(U_1, U_2, Y_1, Y_2 | v_1^n(s_1), v_2^n(s_2)) \right\}.
\end{aligned}$$

Next, consider the following error event in decoding step:

$$\begin{aligned}
\mathbf{E}_8 &= \left(\bigcap_{t=0}^7 \mathbf{E}_t^c \right) \cap \left\{ (V_1^n(\tilde{s}_1), U_1^n(\tilde{s}_1, \tilde{s}'_1), V_2^n(\tilde{s}_2), U_2^n(\tilde{s}_2, \tilde{s}'_2)) \right. \\
&\quad \left. \in \mathcal{T}_\delta^{(n)}(V_1, U_1, V_2, U_2), \right. \\
&\quad \text{for some } (\tilde{s}_1, \tilde{s}'_1, \tilde{s}_2, \tilde{s}'_2) \neq (s_1, s'_1, s_2, s'_2), \\
&\quad \text{s.t. } (V_1^n(\tilde{s}_1), U_1^n(\tilde{s}_1, \tilde{s}'_1), V_2^n(\tilde{s}_2), U_2^n(\tilde{s}_2, \tilde{s}'_2)) \\
&\quad \left. \in \mathcal{B}_1(b_1) \times \mathcal{B}'_1(\tilde{s}_1, w_1) \times \mathcal{B}_2(b_2) \times \mathcal{B}'_2(\tilde{s}_2, w_2) \right\}.
\end{aligned}$$

Finally, by the union of events bound, the probability of error in the encoding and decoding steps is upper bounded as

$$\Pr(\mathbf{E}) \leq \Pr\left(\bigcup_{t=0}^8 \mathbf{E}_t\right) \leq \sum_{t=0}^8 \Pr(\mathbf{E}_t). \tag{54}$$

We proceed to bound each term in (54). From properties of typical sequences, $\Pr(\mathbf{E}_0)$ vanishes as $n \rightarrow \infty$. By covering lemma [39, Lemma 3.3], $\Pr(\mathbf{E}_1)$ and $\Pr(\mathbf{E}_2)$ tend to zero as $n \rightarrow \infty$. For $j, j' = 1, 2$, and $j \neq j'$, since $\{Y_{j'}^n | V_j^n(s_j) = v_j^n, Y_j^n = y_j^n\} \sim \prod_{i=1}^n p(y_{j',i} | y_{j,i})$, by conditional typicality lemma [39, Chapter 2], $\Pr(\mathbf{E}_3)$ tends to zero as $n \rightarrow \infty$. Similarly, as $\{Y_{j'}^n | V_j^n(s_j) = v_j^n, U_j^n(s_j, s'_j) = u_j^n, Y_j^n = y_j^n\} \sim \prod_{i=1}^n p(y_{j',i} | y_{j,i})$, $\Pr(\mathbf{E}_4)$ also vanishes as $n \rightarrow \infty$. To bound $\Pr(\mathbf{E}_5)$, let $(v_{j'}^n, u_{j'}^n, y_1^n, y_2^n) \in \mathcal{T}_\delta^{(n)}(V_{j'}, U_{j'}, Y_1, Y_2)$. Then, $\Pr(V_j^n(s_j) | V_{j'}^n(s_{j'})) = v_j^n, U_j^n(s_j, s'_j) = u_j^n, Y_1^n = y_1^n, Y_2^n = y_2^n = p(v_j^n | y_j^n)$, and by Markov lemma [39,

Lemma 12.1], $\Pr(E_5)$ tends to zero as $n \rightarrow \infty$. Using similar steps and based on Markov lemma, $\Pr(E_6)$ and $\Pr(E_7)$ also tend to zeros as $n \rightarrow \infty$.

As can be seen from E_8 , in the decoding step, an error occurs if the decoded codewords are jointly typical and they are in the bin indexed by (b_1, w_1, b_2, w_2) , however, the decoded tuple $(\hat{s}_1, \hat{s}'_1, \hat{s}_2, \hat{s}'_2)$ of codeword indices are different from the chosen ones at the encoders, i.e., (s_1, s'_1, s_2, s'_2) . We split this event into eight possible events (other events result in the same constraints as one of these eight events) and bound its probability using the union of events bound as follows:

$$\Pr(E_8) = \Pr((\hat{s}_1, \hat{s}'_1, \hat{s}_2, \hat{s}'_2) \neq (s_1, s'_1, s_2, s'_2)) \quad (55)$$

$$\leq \Pr(\hat{s}_1 \neq s_1) \quad (56)$$

$$+ \Pr(\hat{s}_1 = s_1, \hat{s}_2 \neq s_2) \quad (57)$$

$$+ \Pr(\hat{s}_1 \neq s_1, \hat{s}_2 = s_2) \quad (58)$$

$$+ \Pr(\hat{s}_1 = s_1, \hat{s}'_1 \neq s'_1, \hat{s}_2 = s_2) \quad (59)$$

$$+ \Pr(\hat{s}_1 = s_1, \hat{s}'_1 = s'_1, \hat{s}_2 \neq s_2) \quad (60)$$

$$+ \Pr(\hat{s}_1 \neq s_1, \hat{s}_2 = s_2, \hat{s}'_2 = s'_2) \quad (61)$$

$$+ \Pr(\hat{s}_1 = s_1, \hat{s}'_1 = s'_1, \hat{s}_2 = s_2, \hat{s}'_2 \neq s'_2) \quad (62)$$

$$+ \Pr(\hat{s}_1 = s_1, \hat{s}'_1 \neq s'_1, \hat{s}_2 = s_2, \hat{s}'_2 = s'_2). \quad (63)$$

Now, we consider each of the terms in (56)–(63).

$$\begin{aligned} & \Pr(\hat{s}_1 \neq s_1) \\ &= \Pr\left(\exists \hat{s}_1 \neq s_1, \tilde{s}'_1, \tilde{s}_2, \tilde{s}'_2, \text{ s.t. } (v_1^n(\hat{s}_1), u_1^n(\hat{s}_1, \tilde{s}'_1), \right. \\ & \quad \left. v_2^n(\tilde{s}_2), u_2^n(\tilde{s}_2, \tilde{s}'_2)) \in \mathcal{T}_\delta^{(n)}(V_1, U_1, V_2, U_2) \right. \\ & \quad \left. \cap \{\mathcal{B}_1(b_1) \times \mathcal{B}'_1(\hat{s}_1, w_1) \times \mathcal{B}_2(b_2) \times \mathcal{B}'_2(\tilde{s}_2, w_2)\}\right) \\ &\leq \Pr\left((V_1^n, U_1^n, V_2^n, U_2^n) \in \mathcal{T}_\delta^{(n)}(V_1, U_1, V_2, U_2) \right| \\ & \quad (V_1^n, U_1^n) \in \mathcal{T}_\delta^{(n)}(V_1, U_1), (V_2^n, U_2^n) \in \mathcal{T}_\delta^{(n)}(V_2, U_2) \right) \\ & \quad \cdot 2^{n \sum_{j=1}^2 [I(V_j; Y_j) + \epsilon_1 - R_{V_j} + I(U_j; Y_j | V_j) + \epsilon_3 - R_{U_j}]} \\ &\stackrel{(a)}{\leq} 2^{-n(I(V_1, U_1; V_2, U_2) - \eta)} \\ & \quad \cdot 2^{n \sum_{j=1}^2 [I(V_j; Y_j) + \epsilon_1 - R_{V_j} + I(U_j; Y_j | V_j) + \epsilon_3 - R_{U_j}]} \\ &= 2^{n(\sum_{j=1}^2 [I(V_j; Y_j) + \epsilon_1 - R_{V_j} + I(U_j; Y_j | V_j) + \epsilon_3 - R_{U_j}] - I(U_1; U_2) + \eta)}, \end{aligned} \quad (64)$$

where (a) is due to the mutual packing lemma [39, Lemma 12.2]. Therefore, $\Pr(\hat{s}_1 \neq s_1)$ vanishes as $n \rightarrow \infty$ if

$$\begin{aligned} \sum_{j=1}^2 R_{V_j} + R_{U_j} &\geq \sum_{j=1}^2 I(V_j; Y_j) + I(U_j; Y_j | V_j) - I(U_1; U_2) \\ &\stackrel{(a)}{=} I(U_1; Y_1) + I(U_2; Y_2) - I(U_1; U_2) \\ &\stackrel{(b)}{=} I(U_1, U_2; Y_1, Y_2), \end{aligned} \quad (65)$$

where (a) and (b) follow from the long Markov chain $V_1 - U_1 - Y_1 - Y_2 - U_2 - V_2$.

Next, we bound the probability of the event in which s_1 is correctly decoded but not s_2 , i.e., (57) as

$$\begin{aligned} & \Pr(\hat{s}_1 = s_1, \hat{s}_2 \neq s_2) \\ &= \Pr\left(\exists \hat{s}_2 \neq s_2, \tilde{s}'_1, \tilde{s}_2, \text{ s.t. } (v_1^n(s_1), u_1^n(s_1, \tilde{s}'_1), \right. \\ & \quad \left. v_2^n(\tilde{s}_2), u_2^n(\tilde{s}_2, \tilde{s}'_2)) \in \mathcal{T}_\delta^{(n)}(V_1, U_1, V_2, U_2) \right. \\ & \quad \left. \cap \{\mathcal{B}_1(b_1) \times \mathcal{B}'_1(s_1, w_1) \times \mathcal{B}_2(b_2) \times \mathcal{B}'_2(\tilde{s}_2, w_2)\}\right) \\ &\leq \Pr\left((V_1^n, U_1^n, V_2^n, U_2^n) \in \mathcal{T}_\delta^{(n)}(V_1, U_1, V_2, U_2) \right| \\ & \quad (V_1^n, U_1^n) \in \mathcal{T}_\delta^{(n)}(V_1, U_1), (V_2^n, U_2^n) \in \mathcal{T}_\delta^{(n)}(V_2, U_2) \right) \\ & \quad \cdot 2^{n(I(V_2; Y_2) + \epsilon_1 - R_{V_2} + \sum_{j=1}^2 [I(U_j; Y_j | V_j) + \epsilon_3 - R_{U_j}])} \\ &\leq 2^{n(I(V_2; Y_2) + \epsilon_1 - R_{V_2} + \sum_{j=1}^2 [I(U_j; Y_j | V_j) + \epsilon_3 - R_{U_j}] - I(U_1; U_2) + \eta)}, \end{aligned} \quad (66)$$

Therefore, $\Pr(\hat{s}_1 = s_1, \hat{s}_2 \neq s_2)$ vanishes as $n \rightarrow \infty$ if

$$\begin{aligned} R_{V_2} + \sum_{j=1}^2 R_{U_j} &\geq I(V_2; Y_2) + \sum_{j=1}^2 I(U_j; Y_j | V_j) - I(U_1; U_2) \\ &\stackrel{(a)}{=} I(U_1; Y_1 | V_1) + I(U_2; Y_2 | U_1) \\ &\stackrel{(b)}{=} I(U_1; Y_1, Y_2 | V_1) + I(U_2; Y_1, Y_2 | V_1, U_1) \\ &\stackrel{(c)}{=} I(U_1, U_2; Y_1, Y_2 | V_1), \end{aligned} \quad (67)$$

where (a), (b), and (c) follow from the Markov chain $V_1 - U_1 - Y_1 - Y_2 - U_2 - V_2$. With similar steps, $\Pr(\hat{s}_1 \neq s_1, \hat{s}_2 = s_2)$ in (58) tends to zero as $n \rightarrow \infty$ if

$$R_{V_1} + \sum_{j=1}^2 R_{U_j} \geq I(U_1, U_2; Y_1, Y_2 | V_2). \quad (68)$$

We now bound (59) where s_1 and s_2 are decoded correctly:

$$\begin{aligned} & \Pr(\hat{s}_1 = s_1, \hat{s}'_1 \neq s'_1, \hat{s}_2 = s_2) \\ &= \Pr\left(\exists \hat{s}'_1 \neq s'_1, \tilde{s}_2, \text{ s.t. } (v_1^n(s_1), u_1^n(s_1, \hat{s}'_1), \right. \\ & \quad \left. v_2^n(s_2), u_2^n(s_2, \tilde{s}_2)) \in \mathcal{T}_\delta^{(n)}(V_1, U_1, V_2, U_2) \right. \\ & \quad \left. \cap \{\mathcal{B}_1(b_1) \times \mathcal{B}'_1(s_1, w_1) \times \mathcal{B}_2(b_2) \times \mathcal{B}'_2(s_2, w_2)\}\right) \\ &= \Pr\left((V_1^n, U_1^n, V_2^n, U_2^n) \in \mathcal{T}_\delta^{(n)}(V_1, U_1, V_2, U_2) \right| \\ & \quad (V_1^n, U_1^n) \in \mathcal{T}_\delta^{(n)}(V_1, U_1), (V_2^n, U_2^n) \in \mathcal{T}_\delta^{(n)}(V_2, U_2) \right) \\ & \quad \cap \{\mathcal{B}_1(b_1) \times \mathcal{B}'_1(s_1, w_1) \times \mathcal{B}_2(b_2) \times \mathcal{B}'_2(s_2, w_2)\} \\ &= \Pr\left((V_1^n, U_1^n, V_2^n, U_2^n) \in \mathcal{T}_\delta^{(n)}(V_1, U_1, V_2, U_2) \right| \\ & \quad (V_1^n, U_1^n) \in \mathcal{T}_\delta^{(n)}(V_1, U_1), (V_2^n, U_2^n) \in \mathcal{T}_\delta^{(n)}(V_2, U_2) \right) \\ & \quad \cdot \Pr\left((V_1^n, U_1^n, V_2^n, U_2^n) \in \mathcal{T}_\delta^{(n)}(V_1, U_1, V_2, U_2) \right| \\ & \quad (V_1^n, U_1^n) \in \mathcal{T}_\delta^{(n)}(V_1, U_1), (V_1^n, V_2^n) \in \mathcal{T}_\delta^{(n)}(V_1, V_2) \\ & \quad \cap \{\mathcal{B}_1(b_1) \times \mathcal{B}'_1(s_1, w_1) \times \mathcal{B}_2(b_2) \times \mathcal{B}'_2(s_2, w_2)\} \\ &\stackrel{(a)}{\leq} 2^{-n(I(V_1, U_1; V_2 | V_2) - \eta)} \cdot 2^{-n(I(U_1; V_2 | V_1) - \eta)} \\ & \quad \cdot 2^{n \sum_{j=1}^2 [I(U_j; Y_j | V_j) + \epsilon_3 - R_{U_j}]} \end{aligned}$$

$$= 2^{n(\sum_{j=1}^2 [I(U_j; Y_j|V_j) + \epsilon_3 - R_{U_j}] - I(V_1, U_1; U_2|V_2) - I(U_1; V_2|V_1) + 2\eta)}, \quad (69)$$

where (a) follows from the joint typicality lemma and the fact that $U_2^n \sim \prod_{i=1}^n p(u_{2,i}|v_{2,i})$, and $U_1^n \sim \prod_{i=1}^n p(u_{1,i}|v_{1,i})$. Therefore, $\Pr(\hat{s}_1 = s_1, \hat{s}'_1 \neq s'_1, \hat{s}_2 = s_2)$ vanishes as $n \rightarrow \infty$ if

$$\begin{aligned} \sum_{j=1}^2 R_{U_j} &\geq I(U_1; Y_1|V_1) - I(U_1; V_2|V_1) + I(U_2; Y_2|V_2) \\ &\quad - I(V_1, U_1; U_2|V_2) \\ &\stackrel{(a)}{=} I(U_1; Y_1|V_1, V_2) + I(U_2; Y_2|V_1, U_1, V_2) \\ &\stackrel{(b)}{=} I(U_1; Y_1, Y_2|V_1, V_2) + I(U_2; Y_1, Y_2|V_1, U_1, V_2) \\ &= I(U_1, U_2; Y_1, Y_2|V_1, V_2), \end{aligned} \quad (70)$$

where (a) is due the Markov chains $U_1 - Y_1 - V_2$ and $U_2 - Y_1 - (U_1, V_1)$, (b) is due to $U_j - Y_j - Y_{j'}$.

Next, for (60), we have

$$\begin{aligned} \Pr(\hat{s}_1 = s_1, \hat{s}'_1 = s'_1, \hat{s}_2 \neq s_2) &= \Pr\left(\exists \hat{s}_2 \neq s_2, \hat{s}'_2, \text{ s.t. } (v_1^n(s_1), u_1^n(s_1, s'_1), \right. \\ &\quad \left. v_2^n(\hat{s}_2), u_2^n(\hat{s}_2, \hat{s}'_2)) \in \mathcal{T}_\delta^{(n)}(V_1, U_1, V_2, U_2) \right. \\ &\quad \left. \cap \{\mathcal{B}_1(b_1) \times \mathcal{B}'_1(s_1, w_1) \times \mathcal{B}_2(b_2) \times \mathcal{B}'_2(\hat{s}_2, w_2)\}\right) \\ &\leq \Pr\left((V_1^n, U_1^n, V_2^n, U_2^n) \in \mathcal{T}_\delta^{(n)}(V_1, U_1, V_2, U_2) \mid \right. \\ &\quad \left. (V_1^n, U_1^n) \in \mathcal{T}_\delta^{(n)}(V_1, U_1), (V_2^n, U_2^n) \in \mathcal{T}_\delta^{(n)}(V_2, U_2) \right. \\ &\quad \left. \cdot 2^{n(I(V_2; Y_2) + \epsilon_1 - R_{V_2} + I(U_2; Y_2|V_2) + \epsilon_3 - R_{U_2})} \right. \\ &\leq 2^{n(I(V_2; Y_2) + \epsilon_1 - R_{V_2} + I(U_2; Y_2|V_2) + \epsilon_3 - R_{U_2} - I(U_1; U_2) + \eta)}. \end{aligned} \quad (71)$$

Therefore, $\Pr(\hat{s}_1 = s_1, \hat{s}'_1 = s'_1, \hat{s}_2 \neq s_2)$ vanishes as $n \rightarrow \infty$ if

$$\begin{aligned} R_{V_2} + R_{U_2} &\geq I(V_2; Y_2) + I(U_2; Y_2|V_2) - I(U_1; U_2) \\ &= I(U_2; Y_2|U_1), \end{aligned} \quad (72)$$

which follows from the Markov chain $V_2 - U_2 - Y_2 - U_1$. Similarly, $\Pr(\hat{s}_1 \neq s_1, \hat{s}_2 = s_2, \hat{s}'_2 = s'_2)$ in (61) tends to zero as $n \rightarrow \infty$ if

$$R_{V_1} + R_{U_1} \geq I(U_1; Y_1|U_2). \quad (73)$$

We proceed to bound (62) where (s_1, s'_1, s_2) are decoded correctly:

$$\begin{aligned} \Pr(\hat{s}_1 = s_1, \hat{s}'_1 = s'_1, \hat{s}_2 = s_2, \hat{s}'_2 = s'_2) &= \Pr\left(\exists \hat{s}'_2 \neq s'_2, \text{ s.t. } (v_1^n(s_1), u_1^n(s_1, s'_1), \right. \\ &\quad \left. v_2^n(s_2), u_2^n(s_2, \text{hats}'_2)) \in \mathcal{T}_\delta^{(n)}(V_1, U_1, V_2, U_2) \right. \\ &\quad \left. \cap \{\mathcal{B}_1(b_1) \times \mathcal{B}'_1(s_1, w_1) \times \mathcal{B}_2(b_2) \times \mathcal{B}'_2(s_2, w_2)\}\right) \\ &\leq \Pr\left((V_1^n, U_1^n, V_2^n, U_2^n) \in \mathcal{T}_\delta^{(n)}(V_1, U_1, V_2, U_2) \mid \right. \\ &\quad \left. ((V_1^n, U_1^n), V_2^n) \in \mathcal{T}_\delta^{(n)}(V_1, U_1, V_2), \right. \\ &\quad \left. (V_2^n, U_2^n) \in \mathcal{T}_\delta^{(n)}(V_2, U_2) \right. \\ &\quad \left. \cdot 2^{n(I(U_2; Y_2|V_2) + \epsilon_3 - R_{U_2})} \right) \end{aligned}$$

$$\leq 2^{n(I(U_2; Y_2|V_2) + \epsilon_3 - R_{U_2} - I(V_1, U_1; U_2|V_2) + \eta)}, \quad (74)$$

from the joint typicality lemma by considering (V_1, U_1) as one variable and the fact that $U_2^n \sim \prod_{i=1}^n p(u_{2,i}|v_{2,i})$. Therefore, $\Pr(\hat{s}_1 = s_1, \hat{s}'_1 \neq s'_1, \hat{s}_2 = s_2, \hat{s}'_2 = s'_2)$ tends to zero as $n \rightarrow \infty$ if

$$\begin{aligned} R_{U_2} &\geq I(U_2; Y_2|V_2) - I(V_1, U_1; U_2|V_2) \\ &= I(U_2; Y_2|V_2, U_1), \end{aligned} \quad (75)$$

where the last equality is due to $V_1 - U_1 - Y_2 - U_2$. Similarly, $\Pr(\hat{s}_1 = s_1, \hat{s}'_1 \neq s'_1, \hat{s}_2 = s_2, \hat{s}'_2 = s'_2)$ in (63) vanishes as $n \rightarrow \infty$ if

$$R_{U_1} \geq I(U_1; Y_1|V_1, U_2). \quad (76)$$

Gathering (65)–(76), probability of error event E_8 in (55) vanishes as $n \rightarrow \infty$ if, for $j, j' = 1, 2$ and $j \neq j'$, the rates $R_j = R_{V_j} + R_{U_j}$ satisfy

$$R_1 \geq I(U_1; Y_1|U_2), \quad (77)$$

$$R_2 \geq I(U_2; Y_2|U_1), \quad (78)$$

$$R_1 + R_2 \geq I(U_1, U_2; Y_1, Y_2), \quad (79)$$

$$R_{U_j} \geq I(U_j; Y_j|V_j, U_{j'}), \quad (80)$$

$$R_{U_1} + R_{U_2} \geq I(U_1, U_2; Y_1, Y_2|V_1, V_2), \quad (81)$$

$$R_j + R_{U_{j'}} \geq I(U_1, U_2; Y_1, Y_2|V_{j'}). \quad (82)$$

Finally, $\Pr(E)$ in (54) tends to zero as $n \rightarrow \infty$ provided that (77)–(82) are satisfied.

Equivocation rates: The equivocation rate with respect to Agent 1, averaged over all codebooks \mathcal{C}_n , is written as

$$\begin{aligned} H(X^n|f_1(Y_1^n), E^n, \mathcal{C}_n) &= H(X^n|b_1, w_1, E^n, \mathcal{C}_n) \\ &= H(X^n, Y_1^n|b_1, w_1, E^n, \mathcal{C}_n) - H(Y_1^n|b_1, w_1, E^n, X^n, \mathcal{C}_n) \\ &= H(Y_1^n|b_1, w_1, E^n, \mathcal{C}_n) + H(X^n|b_1, w_1, E^n, Y_1^n, \mathcal{C}_n) \\ &\quad - H(Y_1^n|b_1, w_1, E^n, X^n, \mathcal{C}_n) \\ &= H(Y_1^n|b_1, w_1, E^n, \mathcal{C}_n) + H(X^n|E^n, Y_1^n) \\ &\quad - H(Y_1^n|b_1, w_1, X^n, \mathcal{C}_n) \end{aligned} \quad (83)$$

where the last equality holds since $f_1(Y_1^n) = (b_1, w_1)$ is a deterministic function, $Y_1^n - (b_1, w_1, X^n, \mathcal{C}_n) - E^n$ is a Markov chain, and (X^n, Y_1^n, E^n) are independent of the codebook. The first term in (83) is bounded as

$$\begin{aligned} H(Y_1^n|b_1, w_1, E^n, \mathcal{C}_n) &= H(Y_1^n|b_1, E^n, \mathcal{C}_n) - I(Y_1^n; w_1|b_1, E^n, \mathcal{C}_n) \\ &\stackrel{(a)}{\geq} H(Y_1^n|b_1, E^n, \mathcal{C}_n) - H(w_1|b_1, E^n, \mathcal{C}_n) \\ &\stackrel{(b)}{\geq} H(Y_1^n|s_1, E^n, \mathcal{C}_n) - H(w_1|\mathcal{C}_n) \\ &= H(Y_1^n, E^n|s_1, \mathcal{C}_n) - H(E^n|s_1, \mathcal{C}_n) - H(w_1|\mathcal{C}_n) \\ &\stackrel{(c)}{=} H(Y_1^n, E^n) - H(s_1|\mathcal{C}_n) + H(s_1|Y_1^n, E^n, \mathcal{C}_n) \\ &\quad - H(E^n|s_1, \mathcal{C}_n) - H(w_1|\mathcal{C}_n) \\ &\stackrel{(d)}{\geq} H(Y_1^n, E^n) - H(s_1|\mathcal{C}_n) - H(E^n|s_1, \mathcal{C}_n) - H(w_1|\mathcal{C}_n) \\ &\stackrel{(e)}{\geq} nH(Y_1, E) - nI(V_1; Y_1) - n\epsilon_1 - H(E^n|s_1, \mathcal{C}_n) \\ &\quad - nR_{U_1} - n\epsilon_4 \end{aligned}$$

$$\begin{aligned}
&\stackrel{(f)}{\geq} n[H(Y_1, E) - I(V_1; Y_1) - H(E|V_1) - R_{U_1} - \epsilon' \\
&\quad - \epsilon_1 - \epsilon_4] \\
&\stackrel{(g)}{\geq} n[H(Y_1|E) - I(V_1; Y_1|E) - R_{U_1} - \epsilon' - \epsilon_1 - \epsilon_4] \\
&= n[H(Y_1|V_1, E) - R_{U_1} - \epsilon' - \epsilon_1 - \epsilon_4], \tag{84}
\end{aligned}$$

where

- (a) follows as the bin index w_1 is a deterministic function of Y_1^n ;
- (b) follows since the bin index b_1 is a deterministic function of the codeword index s_1 , and conditioning reduces entropy;
- (c) follows since (Y_1^n, E^n) are independent of the codebook;
- (d) follows since the codeword index s_1 is a deterministic function of Y_1^n ;
- (e) follows since random variables $Y_{1,i}$ and E_i are i.i.d., s_1 and w_1 are random variables over sets of size $2^{n(I(V_1; Y_1) + \epsilon_1)}$ and $2^{n(R_{U_1} + \epsilon_4)}$, respectively;
- (f) follows from (52) in Lemma 1;
- (g) follows from the Markov chain $V_1 - Y_1 - E$.

The last term in (83) can be bounded as

$$\begin{aligned}
H(Y_1^n|b_1, w_1, X^n, \mathcal{C}_n) &\stackrel{(a)}{=} H(Y_1^n|b_1, w_1, U_1^n, X^n, \mathcal{C}_n) \\
&= H(Y_1^n|b_1, w_1, U_1^n, U_2^n, X^n, \mathcal{C}_n) \\
&\quad + I(U_1^n; Y_1^n|b_1, w_1, U_2^n, X^n, \mathcal{C}_n) \\
&\stackrel{(b)}{\leq} H(Y_1^n|U_1^n, X^n, \mathcal{C}_n) + H(U_1^n|b_1, w_1, U_2^n, X^n, \mathcal{C}_n) \\
&\stackrel{(c)}{\leq} n[H(Y_1|U_1, X) + \epsilon], \tag{85}
\end{aligned}$$

where

- (a) follows from the Markov chain $Y_1^n - (b_1, w_1, X^n, \mathcal{C}_n) - U_2^n$;
- (b) follows since conditioning reduces entropy, the bin index w_1 is a deterministic function of U_1^n , and the Markov chain $U_1^n - (X^n, \mathcal{C}_n) - U_2^n$;
- (c) follows from Fano's inequality (i.e., in the decoding scheme having U_2^n , b_1 , and w_1 , the CEO decodes U_1^n with high probability).

Substituting (84) and (85) in (83), we obtain

$$\begin{aligned}
H(X^n|f_1(Y_1^n), E^n, \mathcal{C}_n) \\
\geq n[H(Y_1|V_1, E) - R_{U_1} + H(X|Y_1, E) - H(Y_1|U_1, X) \\
- \epsilon - \epsilon' - \epsilon_1 - \epsilon_4]. \tag{86}
\end{aligned}$$

Therefore,

$$\Delta_1 \leq H(Y_1|V_1, E) - R_{U_1} + H(X|Y_1, E) - H(Y_1|U_1, X) \tag{87}$$

is achievable. Based on (80), using Fourier-Motzkin elimination yields

$$\begin{aligned}
\Delta_1 &\leq H(Y_1|V_1, E) - I(U_1; Y_1|V_1, U_2) + H(X|Y_1, E) \\
&\quad - H(Y_1|U_1, X) \\
&\stackrel{(a)}{=} H(Y_1|V_1, E) - I(U_1; Y_1|V_1, U_2) + H(X|Y_1, E, V_1) \\
&\quad - H(Y_1|U_1, V_1, X) \\
&= H(Y_1, X|V_1, E) - I(U_1; Y_1|V_1, U_2) \\
&\quad - H(Y_1|U_1, V_1, X) \\
&= H(X|V_1, E) + H(Y_1|V_1, X, E) - I(U_1; Y_1|V_1, U_2)
\end{aligned}$$

$$\begin{aligned}
&\quad - H(Y_1|U_1, V_1, X) \\
&\stackrel{(b)}{=} H(X|V_1, E) - I(U_1; Y_1|V_1, U_2) + I(U_1; Y_1|V_1, X), \tag{88}
\end{aligned}$$

where

- (a) follows from the Markov chains $V_1 - (U_1, X) - Y_1$ and $V_1 - (Y_1, E) - X$;
- (b) follows from the Markov chain $Y_1 - (V_1, X) - E$.

Similarly for the second link, we obtain

$$\Delta_2 \leq H(X|V_2, E) - I(U_2; Y_2|V_2, U_1) + I(U_2; Y_2|V_2, X). \tag{89}$$

For the sum of the equivocation rates, following the proof of the equivocation rate, we have

$$\begin{aligned}
\Delta_1 + \Delta_2 &\leq H(X|V_1, E) + H(X|V_2, E) - R_{U_1} - R_{U_2} \\
&\quad + I(U_1; Y_1|V_1, X) + I(U_2; Y_2|V_2, X) \\
&\leq H(X|V_1, E) + H(X|V_2, E) \\
&\quad - I(U_1, U_2; Y_1, Y_2|V_1, V_2) \\
&\quad + I(U_1; Y_1|V_1, X) + I(U_2; Y_2|V_2, X), \tag{90}
\end{aligned}$$

where we substitute (81) in the second inequality. In order to prove (8), we have

$$\begin{aligned}
\Delta_1 &\leq H(X|V_1, E) - R_{U_1} + I(U_1; Y_1|V_1, X) \\
&\stackrel{(a)}{\leq} H(X|V_1, E) - I(U_1, U_2; Y_1, Y_2|V_1) + R_2 \\
&\quad + I(U_1; Y_1|V_1, X) \\
&\stackrel{(b)}{=} H(X|V_1, E) - I(U_1; Y_1|V_1) - I(U_2; Y_2|V_1, U_1) \\
&\quad + R_2 + I(U_1; Y_1|V_1, X) \\
&\stackrel{(c)}{=} H(X|V_1, E) - I(U_1; Y_1|V_1) - I(U_2; Y_2|U_1) + R_2 \\
&\quad + I(U_1; Y_1|V_1, X), \tag{91}
\end{aligned}$$

where

- (a) follows from (82);
- (b) follows since $U_1 - Y_1 - Y_2 - U_2$ is a Markov chain;
- (c) follows since $V_1 - U_1 - U_2$ forms a Markov chain.

Inequality (9) can be proved similarly.

Distortion at the CEO: The distortion at the CEO is written as

$$\begin{aligned}
\mathbb{E} \left[d \left(X^n, g(f_1(Y_1^n), f_2(Y_2^n)) \right) \right] &\leq \Pr(\mathsf{E}) d_{\max} \\
&\quad + \Pr(\mathsf{E}^c) \mathbb{E} \left[d \left(X^n, \hat{X}(u_1^n(s_1, s'_1), u_2^n(s_2, s'_2)) \right) \middle| \mathsf{E}^c \right] \\
&\stackrel{(a)}{\leq} \epsilon_n + \frac{1}{n} \sum_{i=1}^n \mathbb{E} \left[d \left(X_i, \hat{X}(u_{1,i}(s_1, s'_1), u_{2,i}(s_2, s'_2)) \right) \middle| \mathsf{E}^c \right] \\
&\stackrel{(b)}{\leq} \epsilon_n + \mathbb{E} \left[d(X, \hat{X}(U_1, U_2)) \right], \tag{92}
\end{aligned}$$

where

- (a) follows since the probability of an error event becomes small as $n \rightarrow \infty$ if the rate constraints are satisfied.

(b) follows since for every jointly typical sequences (x^n, u_1^n, u_2^n) , we have

$$\begin{aligned}
d(x^n, \hat{X}(u_1^n, u_2^n)) &= \frac{1}{n} \sum_{i=1}^n d(x_i, \hat{X}(u_{1,i}, u_{2,i})) \\
&= \frac{1}{n} \sum_{(x, u_1, u_2) \in \mathcal{X} \times \mathcal{U}_1 \times \mathcal{U}_2} \left[d(x, \hat{X}(u_1, u_2)) \right. \\
&\quad \left. N(x, u_1, u_2 | x^n, u_1^n, u_2^n) \right] \\
&= \mathbb{E} \left[d(X, \hat{X}(U_1, U_2)) \right] \\
&\quad + \frac{1}{n} \sum_{(x, u_1, u_2) \in \mathcal{X} \times \mathcal{U}_1 \times \mathcal{U}_2} \left[d(x, \hat{X}(u_1, u_2)) \right. \\
&\quad \left. \left(\frac{1}{n} N(x, u_1, u_2 | x^n, u_1^n, u_2^n) - p(x, u_1, u_2) \right) \right] \\
&\leq \mathbb{E} \left[d(X, \hat{X}(U_1, U_2)) \right] + d_{\max} |\mathcal{X}| |\mathcal{U}_1| |\mathcal{U}_2| \epsilon'_n, \quad (93)
\end{aligned}$$

with $\epsilon'_n \rightarrow 0$ as $n \rightarrow \infty$ from properties of jointly (strong) typical sequences. Therefore, to satisfy the distortion constraint at the CEO, it is sufficient to have

$$\mathbb{E} \left[d(X, \hat{X}(U_1, U_2)) \right] \leq D. \quad (94)$$

This concludes the proof of Theorem 1. \blacksquare

APPENDIX B

PROOF OF PROPOSITION 1: CARDINALITIES IN THE INNER BOUND

To bound the cardinality of alphabets of auxiliary random variables V_1 and U_1 in Theorem 1, we rewrite the equations of the inner bound using straightforward derivations as

$$R_1 \geq H(Y_1|U_2) - H(Y_1|U_1, U_2), \quad (95)$$

$$R_2 \geq I(U_2; Y_2|U_1), \quad (96)$$

$$R_1 + R_2 \geq H(Y_1) - H(Y_1|U_1) + I(U_2; Y_2|U_1), \quad (97)$$

$$\Delta_1 \leq [I(X; U_2|V_1) - I(X; E|V_1) + H(X|U_1, U_2)]^+, \quad (98)$$

$$\Delta_2 \leq [H(X|V_2, E) - H(X|V_2, U_1) + H(X|U_1, U_2)]^+, \quad (99)$$

$$\begin{aligned}
\Delta_1 + \Delta_2 &\leq [I(X; V_2|V_1) - I(X; E|V_1) \\
&\quad + H(X|V_2, E) + H(X|U_1, U_2)]^+, \quad (100)
\end{aligned}$$

$$\begin{aligned}
\Delta_1 - R_2 &\leq [I(X; U_2|V_1) - I(X; E|V_1) \\
&\quad - I(U_2; Y_2|V_1) + H(X|U_1, U_2)]^+, \quad (101)
\end{aligned}$$

$$\begin{aligned}
\Delta_2 - R_1 &\leq [H(X|V_2, E) - H(Y_1|U_2) \\
&\quad + H(Y_1|U_1, U_2) - I(U_2; X|V_2)]^+, \quad (102)
\end{aligned}$$

$$D \geq \mathbb{E} \left[d(X, \hat{X}(U_1, U_2)) \right]. \quad (103)$$

Using standard arguments based on the Fenchel-Eggleston-Carathéodory theorem and the support lemma [39, Appendix C], it can be shown that \mathcal{V}_1 should have $|\mathcal{Y}_1| - 1$ letters to preserve the probability distribution $p(y_1)$, and ten more to preserve $H(Y_1|U_1, U_2)$, $I(U_2; Y_2|U_1)$, $H(Y_1|U_1)$, $I(X; U_2|V_1)$, $I(X; E|V_1)$, $H(X|U_1, U_2)$, $H(X|V_2, U_1)$, $I(X; V_2|V_1)$, $I(U_2; Y_2|V_1)$, and the distortion constraint.

Furthermore, for each $v_1 \in \mathcal{V}_1$, there exists a set \mathcal{U}'_1 with $|\mathcal{Y}_1| - 1$ letters to preserve the distribution $p(y_1)$, plus six more to preserve $H(Y_1|U_1, U_2)$, $I(U_2; Y_2|U_1)$, $H(Y_1|U_1)$, $H(X|U_1, U_2)$, $H(X|V_2, U_1)$, and the distortion constraint.

Therefore, it suffices to have $|\mathcal{V}_1| \leq |\mathcal{Y}_1| + 9$ and $|\mathcal{U}_1| \leq (|\mathcal{Y}_1| + 9)(|\mathcal{Y}_1| + 5)$. The same holds for the cardinalities of the sets \mathcal{V}_2 and \mathcal{U}_2 . \blacksquare

APPENDIX C PROOF OF THEOREM 2: THE OUTER BOUND

We denote the message transmitted by the j th agent as $J_j = f_j(Y_j^n)$, and also define auxiliary random variables

$$V_{j,i} := (J_j, X^{i-1}), \quad (104)$$

$$\begin{aligned}
U_{j,i} &:= (V_{j,i}, Y_1^{i-1}, Y_2^{i-1}) \\
&= (J_j, X^{i-1}, Y_1^{i-1}, Y_2^{i-1}),
\end{aligned} \quad (105)$$

for $j \in \{1, 2\}$ and $i \in \{1, \dots, n\}$. Then, we have the following chain of inequalities for Agent 1's rate:

$$\begin{aligned}
n(R_1 + \epsilon) &\geq H(J_1) \\
&\stackrel{(a)}{=} I(J_1; J_2, X^n, Y_1^n, Y_2^n) \\
&\stackrel{(b)}{\geq} I(J_1; X^n, Y_1^n, Y_2^n | J_2) \\
&\stackrel{(c)}{=} \sum_{i=1}^n I(J_1; X_i, Y_{1,i}, Y_{2,i} | J_2, X^{i-1}, Y_1^{i-1}, Y_2^{i-1}) \\
&\stackrel{(d)}{\geq} \sum_{i=1}^n I(J_1; Y_{1,i} | J_2, X^{i-1}, Y_1^{i-1}, Y_2^{i-1}) \\
&= \sum_{i=1}^n I(J_1, X^{i-1}, Y_1^{i-1}, Y_2^{i-1}; Y_{1,i} | J_2, X^{i-1}, Y_1^{i-1}, Y_2^{i-1}) \\
&\stackrel{(e)}{=} \sum_{i=1}^n I(U_{1,i}; Y_{1,i} | U_{2,i}),
\end{aligned} \quad (106)$$

where

- (a) follows as $J_1 = f_1(Y_1^n)$;
- (b) follows from the non-negativity of mutual information;
- (c) follows from the chain rule of conditional mutual information;
- (d) follows from the non-negativity of mutual information;
- (e) follows from the definition of auxiliary random variables in (104)–(105).

The equivocation with respect to Agent 1 is bounded as

$$\begin{aligned}
n(\Delta_1 - \epsilon) &\leq H(X^n | J_1, E^n) \\
&= H(X^n | J_1) - I(X^n; E^n | J_1) \\
&\stackrel{(a)}{=} H(X^n | J_1) - I(X^n; E^n) + I(J_1; E^n) \\
&\stackrel{(b)}{=} \sum_{i=1}^n \left[H(X_i | J_1, X^{i-1}) - I(X_i; E_i) + I(J_1, E^{i-1}; E_i) \right] \\
&\stackrel{(c)}{\leq} \sum_{i=1}^n \left[H(X_i | J_1, X^{i-1}) + I(J_1, X^{i-1}, E^{i-1}; E_i) \right. \\
&\quad \left. - I(X_i; E_i) \right] \\
&\stackrel{(d)}{=} \sum_{i=1}^n \left[H(X_i | J_1, X^{i-1}) - I(X_i; E_i | J_1, X^{i-1}) \right]
\end{aligned}$$

$$\begin{aligned}
&\stackrel{(e)}{=} \sum_{i=1}^n \left[H(X_i|V_{1,i}) - I(X_i; E_i|V_{1,i}) \right] \\
&= \sum_{i=1}^n \left[H(X_i|E_i) - I(X_i; V_{1,i}|E_i) \right]
\end{aligned} \tag{107}$$

where

- (a) follows since $J_1 - X^n - E^n$ is a Markov chain;
- (b) follows since X_i and E_i are memoryless, and the chain rule of conditional entropy;
- (c) follows the non-negativity of mutual information;
- (d) follows from the Markov chain $E_i - X_i - (J_1, X^{i-1}) - E^{i-1}$;
- (e) follows from the definition of auxiliary random variable in (104).

To prove (16), we have

$$\begin{aligned}
n(\Delta_1 - \epsilon) &\leq H(X^n|J_1, E^n) \\
&= H(X^n|J_1, J_2, E^n) + I(X^n; J_2|J_1, E^n) \\
&\stackrel{(a)}{\leq} \sum_{i=1}^n \left[H(X_i|J_1, J_2, X^{i-1}, E^n) \right] + H(J_2) \\
&\stackrel{(b)}{\leq} \sum_{i=1}^n \left[H(X_i|J_1, J_2, X^{i-1}, E_i) \right] + H(J_2) \\
&\stackrel{(c)}{=} \sum_{i=1}^n \left[H(X_i|V_{1,i}, V_{2,i}, E_i) \right] + H(J_2) \\
&\leq \sum_{i=1}^n \left[H(X_i|V_{1,i}, E_i) - I(X_i; V_{2,i}|V_{1,i}, E_i) \right] + nR_2 + n\epsilon,
\end{aligned} \tag{108}$$

where

- (a) follows from the non-negativity of conditional mutual information and the fact that conditioning reduces entropy;
- (b) follows since conditioning reduces entropy;
- (c) follows from definition of auxiliary random variables (104).

Since the setup is symmetric, bounds on Agent 2's transmission and equivocation rates in (12), (15), and (17) are derived similar to those of Agent 1. The sum rate is bounded as

$$\begin{aligned}
n(R_1 + R_2 + 2\epsilon) &\geq H(J_1, J_2) \\
&\stackrel{(a)}{=} I(J_1, J_2; X^n, Y_1^n, Y_2^n) \\
&\stackrel{(b)}{=} \sum_{i=1}^n I(J_1, J_2; X_i, Y_{1,i}, Y_{2,i} | X^{i-1}, Y_1^{i-1}, Y_2^{i-1}) \\
&= \sum_{i=1}^n I(J_1, J_2, X^{i-1}, Y_1^{i-1}, Y_2^{i-1}; X_i, Y_{1,i}, Y_{2,i}) \\
&\quad - I(X^{i-1}, Y_1^{i-1}, Y_2^{i-1}; X_i, Y_{1,i}, Y_{2,i}) \\
&\stackrel{(c)}{=} \sum_{i=1}^n I(J_1, J_2, X^{i-1}, Y_1^{i-1}, Y_2^{i-1}; X_i, Y_{1,i}, Y_{2,i}) \\
&\stackrel{(d)}{\geq} \sum_{i=1}^n I(J_1, J_2, X^{i-1}, Y_1^{i-1}, Y_2^{i-1}; Y_{1,i}, Y_{2,i}) \\
&\stackrel{(e)}{=} \sum_{i=1}^n I(U_{1,i}, U_{2,i}; Y_{1,i}, Y_{2,i}),
\end{aligned} \tag{109}$$

where

- (a) follows as $J_1 = f_1(Y_1^n)$ and $J_2 = f_2(Y_2^n)$;
- (b) follows from the chain rule of mutual information;
- (c) follows since $X_i, Y_{1,i}, Y_{2,i}$ are memoryless;
- (d) follows from the non-negativity of mutual information;
- (e) follows from definition of auxiliary random variables (105).

Next, the distortion bound at the CEO is written as

$$\begin{aligned}
D + \epsilon &\geq \mathbb{E} \left[d(X^n, g(J_1, J_2)) \right] \\
&\geq \frac{1}{n} \sum_{i=1}^n \mathbb{E} \left[d(X_i, \hat{X}_i(U_{1,i}, U_{2,i})) \right],
\end{aligned} \tag{110}$$

where $\hat{X}_i(U_{1,i}, U_{2,i}) := g_i(J_1, J_2)$ is the i th element of the CEO's decoded sequence.

We define a time sharing random variable Q independent from all other random variables and uniformly distributed over the set $\{1, \dots, n\}$. We further define $X = X_Q$, $Y_j = Y_{j,Q}$, $E = E_Q$, $U_j = (Q, U_{j,Q})$, and $V_j = (Q, V_{j,Q})$ for $j \in \{1, 2\}$. Also let $\hat{X}(U_1, U_2) = \hat{X}(Q, U_{1,Q}, U_{2,Q}) = \hat{X}_Q(U_{1,Q}, U_{2,Q})$. We have the Markov chains $V_j - U_j - Y_j - (X, E, Y_{j'})$ for $j, j' \in \{1, 2\}$ and $j \neq j'$. Now, we can write the rates and equivocation bounds for Agent 1 (and similarly for Agent 2) as follows:

$$\begin{aligned}
R_1 + \epsilon &\geq \frac{1}{n} \sum_{i=1}^n I(U_{1,i}; Y_{1,i}|U_{2,i}) \\
&= \frac{1}{n} \sum_{i=1}^n I(U_{1,Q}; Y_{1,Q}|U_{2,Q}, Q = i) \\
&= I(U_1; Y_1|U_2),
\end{aligned} \tag{111}$$

$$\begin{aligned}
\Delta_1 - \epsilon &\leq \frac{1}{n} \sum_{i=1}^n \left[H(X_i|E_i) - I(X_i; V_{1,i}|E_i) \right] \\
&= \frac{1}{n} \sum_{i=1}^n \left[H(X_Q|E_Q, Q = i) \right. \\
&\quad \left. - I(X_Q; V_{1,Q}|E_Q, Q = i) \right] \\
&= H(X|E) - I(X; V_1|E),
\end{aligned} \tag{112}$$

$$\begin{aligned}
\Delta_1 - R_2 - 2\epsilon &\leq \frac{1}{n} \sum_{i=1}^n \left[H(X_i|V_{1,i}, E_i) - I(X_i; V_{2,i}|V_{1,i}, E_i) \right] \\
&= \frac{1}{n} \sum_{i=1}^n \left[H(X_Q|V_{1,Q}, E_Q, Q = i) \right. \\
&\quad \left. - I(X_Q; V_{2,Q}|V_{1,Q}, E_Q, Q = i) \right] \\
&= H(X|V_1, E) - I(X; V_2|V_1, E).
\end{aligned} \tag{113}$$

The sum rate and the distortion bounds can also be written as

$$\begin{aligned}
R_1 + R_2 + 2\epsilon &\geq \frac{1}{n} \sum_{i=1}^n I(U_{1,i}, U_{2,i}; Y_{1,i}, Y_{2,i}) \\
&= \frac{1}{n} \sum_{i=1}^n I(U_{1,Q}, U_{2,Q}; Y_{1,Q}, Y_{2,Q}|Q = i) \\
&= I(U_1, U_2; Y_1, Y_2),
\end{aligned} \tag{114}$$

$$\begin{aligned}
D + \epsilon &\geq \frac{1}{n} \sum_{i=1}^n \mathbb{E} \left[d \left(X_i, \hat{X}_i(U_{1,i}, U_{2,i}) \right) \right] \\
&= \frac{1}{n} \sum_{i=1}^n \mathbb{E} \left[d \left(X_Q, \hat{X}_Q(U_{1,Q}, U_{2,Q}) \right) \middle| Q = i \right] \\
&= \mathbb{E} \left[d \left(X_Q, \hat{X}_Q(U_{1,Q}, U_{2,Q}) \right) \right] \\
&= \mathbb{E} \left[d \left(X, \hat{X}(U_1, U_2) \right) \right]. \tag{115}
\end{aligned}$$

Letting ϵ and ϵ_n tend to zero, inequalities of the Theorem 2 follow. \blacksquare

APPENDIX D

PROOF OF PROPOSITION 2: CARDINALITIES IN THE OUTER BOUND

To bound the cardinality of the sets \mathcal{V}_1 and \mathcal{U}_1 in Theorem 2, we rewrite the equations of the outer bound as

$$R_1 \geq H(Y_1|U_2) - H(Y_1|U_1, U_2), \tag{116}$$

$$R_2 \geq I(U_2; Y_2|U_1), \tag{117}$$

$$R_1 + R_2 \geq H(Y_1) - H(Y_1|U_1) + I(U_2; Y_2|U_1), \tag{118}$$

$$\Delta_1 \leq H(X|V_1) - I(X; E|V_1), \tag{119}$$

$$\Delta_2 \leq H(X|V_2) - I(X; E|V_2), \tag{120}$$

$$\begin{aligned}
\Delta_1 - R_2 &\leq H(X|V_1) - I(X; E|V_1) - I(X; V_2|V_1) \\
&\quad + I(E; V_2|V_1), \tag{121}
\end{aligned}$$

$$\begin{aligned}
\Delta_2 - R_1 &\leq H(X|V_1) - I(X; V_2|V_1) - I(X; E|V_1) \\
&\quad + I(E; V_2|V_1), \tag{122}
\end{aligned}$$

$$D \geq \mathbb{E} \left[d \left(X, \hat{X}(U_1, U_2) \right) \right]. \tag{123}$$

Similar to the approach in Appendix B, using standard arguments based on the Fenchel-Eggleston-Carathéodory theorem and the support lemma [39, Appendix C], it can be shown that \mathcal{V}_1 should have $|\mathcal{Y}_1| - 1$ letters to preserve the probability distribution $p(y_1)$, and eight more to preserve $H(Y_1|U_1, U_2)$, $I(U_2; Y_2|U_1)$, $H(Y_1|U_1)$, $H(X|V_1)$, $I(X; E|V_1)$, $I(X; V_2|V_1)$, $I(E; V_2|V_1)$, and the distortion constraint. Furthermore, for each $v_1 \in \mathcal{V}_1$, there exists a set \mathcal{U}'_1 with $|\mathcal{Y}_1| - 1$ letters to preserve the distribution $p(y_1)$, plus four more to preserve $H(Y_1|U_1, U_2)$, $I(U_2; Y_2|U_1)$, $H(Y_1|U_1)$, and the distortion constraint.

Therefore, it suffices to have $|\mathcal{V}_1| \leq |\mathcal{Y}_1| + 7$ and $|\mathcal{U}_1| \leq (|\mathcal{Y}_1| + 7)(|\mathcal{Y}_1| + 3)$. The same holds for the cardinalities of the sets \mathcal{V}_2 and \mathcal{U}_2 . \blacksquare

APPENDIX E CONVERSE PROOF OF COROLLARY 4

Let $J_1 = f_1(X^n)$ and $J_2 = f_2(Y_2^n)$ denote the messages transmitted by Agent 1 and Agent 2, respectively, and also define the auxiliary random variable $U_{2,i} = (J_2, X^{i-1})$, for $i \in \{1, \dots, n\}$.

For the equivocation rates, we have

$$\begin{aligned}
n(\Delta_1 - \epsilon) &\leq H(X^n|J_1) \\
&= H(X^n, J_2|J_1) - H(J_2|J_1 X^n) \\
&= H(J_2|J_1) - H(J_2|J_1 X^n) + H(X^n|J_1 J_2) \\
&\stackrel{(a)}{\leq} H(J_2) - H(J_2|X^n) + n\epsilon_n \\
&= I(J_2; X^n) + n\epsilon_n \\
&= \sum_{i=1}^n I(J_2 X^{i-1}; X_i) + n\epsilon_n \\
&= \sum_{i=1}^n I(U_{2,i}; X_i) + n\epsilon_n, \tag{124}
\end{aligned}$$

where (a) follows from Fano's inequality, the Markov chain $J_1 - X^n - J_2$, and the fact that conditioning reduces entropy. We also have

$$\begin{aligned}
n(\Delta_2 - \epsilon) &\leq H(X^n|J_2) \\
&= \sum_{i=1}^n H(X_i|J_2 X^{i-1}) \\
&= \sum_{i=1}^n H(X_i|U_{2,i}). \tag{125}
\end{aligned}$$

Similar to the proof of Theorem 2, by using an independent random variable Q , inequalities of Corollary 4 follow. \blacksquare

APPENDIX F PROOF OF THEOREM 3: QUADRATIC GAUSSIAN CASE WITHOUT SIDE INFORMATION AT EVE

A. Achievability

The achievability is proved by applying the results of Theorem 1 to the Gaussian case. Although Theorem 1 is proved for finite alphabet sources with bounded distortion measure, its results can be applied to the Gaussian sources with quadratic distortion measure as shown e.g., in [40]. In addition similar results are proved for the Gaussian CEO problem (without secrecy constraints) in [13] and [14].

It follows from Theorem 1 that with no side information at Eve the following region is achievable:

$$R_1 \geq I(U_1; Y_1|U_2), \tag{126}$$

$$R_2 \geq I(U_2; Y_2|U_1), \tag{127}$$

$$R_1 + R_2 \geq I(U_1, U_2; Y_1, Y_2), \tag{128}$$

$$\Delta_1 \leq h(X) - I(U_1; Y_1|U_2) + I(U_1; Y_1|X), \tag{129}$$

$$\Delta_2 \leq h(X) - I(U_2; Y_2|U_1) + I(U_2; Y_2|X), \tag{130}$$

$$\begin{aligned}
\Delta_1 + \Delta_2 &\leq 2h(X) - I(U_1, U_2; Y_1, Y_2) \\
&\quad + I(U_1; Y_1|X) + I(U_2; Y_2|X), \tag{131}
\end{aligned}$$

$$\begin{aligned}
\Delta_1 - R_2 &\leq h(X) - I(U_2; Y_2|U_1) - I(U_1; Y_1) \\
&\quad + I(U_1; Y_1|X), \tag{132}
\end{aligned}$$

$$\begin{aligned}
\Delta_2 - R_1 &\leq h(X) - I(U_1; Y_1|U_2) - I(U_2; Y_2) \\
&\quad + I(U_2; Y_2|X), \tag{133}
\end{aligned}$$

$$D \geq \mathbb{E} \left[d \left(X, \hat{X}(U_1, U_2) \right) \right], \tag{134}$$

where $h(\cdot)$ is the differential entropy. This region is obtained by setting V_1 , V_2 , and E to be constant in (5)–(9).

Let U_1 and U_2 be auxiliary random variables jointly distributed with the source X and the respective agent's observation Y_1 and Y_2 such that

$$U_j = Y_j + Z_j, \quad j \in \{1, 2\} \quad (135)$$

where $Z_j \sim \mathcal{N}(0, \sigma_{Z_j}^2)$ for $j \in \{1, 2\}$ is independent of X , Y_1 , and Y_2 . Given U_1 and U_2 , we choose the estimator function $\hat{X}(U_1, U_2)$ to be the minimum mean square error (MMSE) estimator. For a fixed target distortion $D > 0$, the auxiliary random variables are adjusted so that the distortion constraint (134) is satisfied. Therefore, we have

$$D = \mathbb{E}[d(X, \hat{X}(U_1, U_2))] = \mathbb{E}[(X - \hat{X}(U_1, U_2))^2] = \frac{1}{2\pi e} 2^{2h(X|U_1, U_2)}, \quad (136)$$

where (136) follows from [37, Theorem 8.6.6]. We then calculate

$$h(X|U_1, U_2) = h(X, U_1, U_2) - h(U_1, U_2) = \frac{1}{2} \log \left((2\pi e) \frac{\det \mathbf{K}_{XU_1U_2}}{\det \mathbf{K}_{U_1U_2}} \right), \quad (137)$$

where the covariance matrix $\mathbf{K}_{XU_1U_2}$ is

$$\mathbf{K}_{XU_1U_2} = \begin{bmatrix} \sigma_X^2 & \sigma_X^2 & \sigma_X^2 \\ \sigma_X^2 & \sigma_X^2 + \sigma_{N_1}^2 + \sigma_{Z_1}^2 & \sigma_X^2 \\ \sigma_X^2 & \sigma_X^2 & \sigma_X^2 + \sigma_{N_2}^2 + \sigma_{Z_2}^2 \end{bmatrix}.$$

Substituting (137) in (136), we obtain

$$\frac{1}{D} = \frac{1}{\sigma_X^2} + \frac{1}{\sigma_{N_1}^2 + \sigma_{Z_1}^2} + \frac{1}{\sigma_{N_2}^2 + \sigma_{Z_2}^2}. \quad (138)$$

Next, for $j \in \{1, 2\}$, we define

$$r_j := I(U_j; Y_j|X) \quad (139)$$

$$= I(U_j; Y_j) - I(U_j; X) \quad (140)$$

$$= \frac{1}{2} \log \frac{\sigma_{N_j}^2 + \sigma_{Z_j}^2}{\sigma_{Z_j}^2}, \quad (141)$$

where the second equality is due to the Markov chain $U_j - Y_j - X$. For any $r_j \geq 0$ with $j \in \{1, 2\}$, there exists corresponding $\sigma_{Z_j}^2$, and thus, auxiliary random variable U_j . Then, using (141), we can rewrite (138) as

$$\frac{1}{D} = \frac{1}{\sigma_X^2} + \frac{1 - 2^{-2r_1}}{\sigma_{N_1}^2} + \frac{1 - 2^{-2r_2}}{\sigma_{N_2}^2}. \quad (142)$$

Now, for the rate of Agent 1, from (126) we have

$$\begin{aligned} R_1 &\geq I(U_1; Y_1|U_2) \\ &\stackrel{(a)}{=} I(U_1; X, Y_1|U_2) \\ &\stackrel{(b)}{=} I(U_1; X|U_2) + I(U_1; Y_1|X) \\ &\stackrel{(c)}{=} h(X|U_2) - h(X|U_1, U_2) + r_1 \\ &\stackrel{(d)}{=} -\frac{1}{2} \log \left(\frac{1}{2\pi e} \left(\frac{1}{\sigma_X^2} + \frac{1 - 2^{-2r_2}}{\sigma_{N_2}^2} \right) \right) \\ &\quad - \frac{1}{2} \log(2\pi e D) + r_1 \\ &= \frac{1}{2} \log \frac{1}{D} - \frac{1}{2} \log \left(\frac{1}{\sigma_X^2} + \frac{1 - 2^{-2r_2}}{\sigma_{N_2}^2} \right) + r_1, \end{aligned} \quad (143)$$

where

- (a) follows as $U_1 - Y_1 - (X, U_2)$ forms a Markov chain;
- (b) follows as $(U_1, Y_1) - X - U_2$ is also a Markov chain;
- (c) follows from the definition of r_1 in (139);
- (d) follows from (136) and (141).

The rate of Agent 2 in (127) is written in a similar way. For the sum rate, we have

$$\begin{aligned} R_1 + R_2 &\geq I(U_1, U_2; Y_1, Y_2) \\ &\stackrel{(a)}{=} I(U_1, U_2; X, Y_1, Y_2) \\ &\stackrel{(b)}{=} I(U_1, U_2; X) + I(U_1, U_2; Y_1, Y_2|X) \\ &\stackrel{(c)}{=} h(X) - h(X|U_1, U_2) + I(U_1; Y_1|X) + I(U_2; Y_2|X) \\ &\stackrel{(d)}{=} \frac{1}{2} \log(2\pi e \sigma_X^2) - \frac{1}{2} \log(2\pi e D) + r_1 + r_2 \\ &= \frac{1}{2} \log \frac{1}{D} - \frac{1}{2} \log \frac{1}{\sigma_X^2} + r_1 + r_2, \end{aligned} \quad (144)$$

where (a)–(c) are due to the long Markov chain $U_1 - Y_1 - X - Y_2 - U_2$ and (d) from (136) and (139).

Using (139) and the bound on R_1 (and R_2), the equivocation rate in (129) (and similarly (130)) can be written as

$$\begin{aligned} \Delta_1 &\leq h(X) - I(U_1; Y_1|U_2) + I(U_1; Y_1|X) \\ &= \frac{1}{2} \log(2\pi e \sigma_X^2) - \frac{1}{2} \log \frac{1}{D} + \frac{1}{2} \log \left(\frac{1}{\sigma_X^2} + \frac{1 - 2^{-2r_2}}{\sigma_{N_2}^2} \right). \end{aligned} \quad (145)$$

For sum of the equivocation rates, we have

$$\begin{aligned} \Delta_1 + \Delta_2 &\leq 2h(X) - I(U_1, U_2; Y_1, Y_2) + I(U_1; Y_1|X) \\ &\quad + I(U_2; Y_2|X) \\ &= \log(2\pi e \sigma_X^2) - \frac{1}{2} \log \frac{1}{D} + \frac{1}{2} \log \frac{1}{\sigma_X^2} \\ &= \log(2\pi e) - \frac{1}{2} \log \frac{1}{D} - \frac{1}{2} \log \frac{1}{\sigma_X^2}. \end{aligned} \quad (146)$$

Finally, the bound in (132) (and similarly (133)) can be

rewritten as

$$\begin{aligned}
\Delta_1 - R_2 &\leq h(X) - I(U_2; Y_2|U_1) - I(U_1; Y_1) + I(U_1; Y_1|X) \\
&\stackrel{(a)}{=} h(X) - I(U_1, U_2; Y_1, Y_2) + I(U_1; Y_1|X) \\
&= \log(2\pi e\sigma_X^2) - \frac{1}{2} \log \frac{1}{D} + \frac{1}{2} \log \frac{1}{\sigma_X^2} - r_1 - r_2 + r_1 \\
&= \log(2\pi e) - \frac{1}{2} \log \frac{1}{D} - r_2,
\end{aligned} \tag{147}$$

where (a) is due to the Markov chains $U_1 - Y_1 - Y_2$ and $Y_1 - Y_2 - U_2$. \blacksquare

B. Converse

Denote the messages transmitted by the agents as $J_1 = f_1(Y_1^n)$ and $J_2 = f_2(Y_2^n)$, and define

$$r_j := \frac{1}{n} I(J_j; Y_j^n | X^n). \tag{148}$$

Let \mathcal{J} be a subset of $\{1, 2\}$ and \mathcal{J}^c its complement. We have

$$\begin{aligned}
\sum_{j \in \mathcal{J}} n(R_j + \epsilon) &\geq H(J_{\mathcal{J}}) \\
&\geq H(J_{\mathcal{J}} | J_{\mathcal{J}^c}) \\
&= I(J_{\mathcal{J}}; Y_{\mathcal{J}}^n | J_{\mathcal{J}^c}) \\
&\stackrel{(a)}{=} I(J_{\mathcal{J}}; X^n, Y_{\mathcal{J}}^n | J_{\mathcal{J}^c}) \\
&= I(J_{\mathcal{J}}; X^n | J_{\mathcal{J}^c}) + I(J_{\mathcal{J}}; Y_{\mathcal{J}}^n | X^n, J_{\mathcal{J}^c}) \\
&\stackrel{(b)}{=} I(J_{\mathcal{J}}, J_{\mathcal{J}^c}; X^n) - I(J_{\mathcal{J}^c}; X^n) + \sum_{j \in \mathcal{J}} I(J_j; Y_j^n | X^n) \\
&\stackrel{(c)}{\geq} \left[I(\hat{X}^n; X^n) - I(J_{\mathcal{J}^c}; X^n) \right]^+ + \sum_{j \in \mathcal{J}} n r_j, \tag{149}
\end{aligned}$$

where

- (a) follows from the Markov chain $J_{\mathcal{J}} - (Y_{\mathcal{J}}^n, J_{\mathcal{J}^c}) - X^n$;
- (b) follows from the Markov chain $(J_j, Y_j^n) - X^n - J_{j'}$ for $j \neq j'$;
- (c) follows from the data processing inequality and the definition of r_j in (148).

The first term in the right-hand side of (149) is bounded as

$$\begin{aligned}
I(\hat{X}^n; X^n) &= h(X^n) - h(X^n | \hat{X}^n) \\
&= h(X^n) - h(X^n - \hat{X}^n | \hat{X}^n) \\
&\geq h(X^n) - h(X^n - \hat{X}^n) \\
&\geq \frac{n}{2} \log(2\pi e\sigma_X^2) - \frac{n}{2} \log(2\pi eD) \\
&= \frac{n}{2} \log \frac{\sigma_X^2}{D}.
\end{aligned} \tag{150}$$

We use the following lemma to bound the second term in the right-hand side of (149). The proof of the lemma is given in [13] and [14].

Lemma 2: Define $r_j := \frac{1}{n} I(J_j; Y_j^n | X^n)$ for $j \in \{1, 2\}$. Then, for $\mathcal{J} \subseteq \{1, 2\}$,

$$\frac{1}{n} I(J_{\mathcal{J}}; X^n) \leq \frac{1}{2} \log \left[\sigma_X^2 \left(\frac{1}{\sigma_X^2} + \sum_{j \in \mathcal{J}} \frac{1 - 2^{-2r_j}}{\sigma_{N_j}^2} \right) \right].$$

Based on Lemma 2, (150), and (149), we obtain

$$\begin{aligned}
\sum_{j \in \mathcal{J}} R_j &\geq \frac{1}{2} \log \frac{1}{D} - \frac{1}{2} \log \left(\frac{1}{\sigma_X^2} + \sum_{j \in \mathcal{J}^c} \frac{1 - 2^{-2r_j}}{\sigma_{N_j}^2} \right) \\
&\quad + \sum_{j \in \mathcal{J}} r_j.
\end{aligned} \tag{151}$$

Now, substituting \mathcal{J} with sets $\{1\}$, $\{2\}$, $\{1, 2\}$, and \emptyset leads to the bounds (30)–(32) and (38) of Theorem 3, respectively.

The bounds on the equivocation rates are obtained as

$$\begin{aligned}
\sum_{j \in \mathcal{J}} n(\Delta_j - \epsilon) &\leq \sum_{j \in \mathcal{J}} h(X^n | J_j) \\
&= \sum_{j \in \mathcal{J}} h(X^n) - H(J_j) + H(J_j | X^n) \\
&\stackrel{(a)}{=} \sum_{j \in \mathcal{J}} h(X^n) - H(J_j) + I(J_j; Y_j^n | X^n) \\
&\stackrel{(b)}{\leq} \sum_{j \in \mathcal{J}} \left[h(X^n) + I(J_j; Y_j^n | X^n) \right] - H(J_{\mathcal{J}}) \\
&\stackrel{(c)}{\leq} \sum_{j \in \mathcal{J}} \frac{n}{2} \log(2\pi e\sigma_X^2) - \frac{n}{2} \log \frac{1}{D} \\
&\quad + \frac{n}{2} \log \left(\frac{1}{\sigma_X^2} + \sum_{j \in \mathcal{J}^c} \frac{1 - 2^{-2r_j}}{\sigma_{N_j}^2} \right), \tag{152}
\end{aligned}$$

where

- (a) follows since $J_j = f_j(Y_j^n)$;
- (b) follows since $H(J_{\mathcal{J}}) \leq \sum_{j \in \mathcal{J}} H(J_j)$;
- (c) follows from the definition of r_j in (148) and the same techniques used to bound $H(J_{\mathcal{J}})$ in (149).

Setting \mathcal{J} to $\{1\}$, $\{2\}$, and $\{1, 2\}$ leads to the bounds (33)–(35), respectively.

For the bound in (36), consider

$$\begin{aligned}
n(\Delta_1 - \epsilon) &\leq h(X^n | J_1) \\
&= h(X^n | J_1, J_2) + I(X^n; J_2 | J_1) \\
&= h(X^n) - I(X^n; J_1, J_2) + H(J_2 | J_1) - H(J_2 | J_1, X^n) \\
&\stackrel{(a)}{\leq} h(X^n) - I(X^n; \hat{X}^n) + H(J_2) - H(J_2 | X^n) \\
&\stackrel{(b)}{\leq} h(X^n) - I(X^n; \hat{X}^n) + nR_2 - I(J_2; Y_2^n | X^n) \\
&\stackrel{(c)}{\leq} \frac{n}{2} \log(2\pi e\sigma_X^2) - \frac{n}{2} \log \frac{\sigma_X^2}{D} + nR_2 - nr_2, \tag{153}
\end{aligned}$$

where

- (a) follows since $J_1 - X^n - J_2$ and that conditioning reduces entropy;
- (b) follows since $J_2 = f_2(Y_2^n)$;
- (c) follows from (150) and the definition of r_j in (148).

Therefore,

$$\Delta_1 - R_2 \leq \frac{1}{2} \log(2\pi e) - \frac{1}{2} \log \frac{1}{D} - r_2. \tag{154}$$

The bound in (37) can be proved similarly. \blacksquare

APPENDIX G

PROOF OF THEOREM 4: QUADRATIC GAUSSIAN CASE
WITH SIDE INFORMATION AT EVE

The proof of Theorem 4 is based on a similar approach for the achievability proof of Theorem 3 described in Appendix F-A. However, in this case due to the side information at Eve, setting the auxiliary random variable V_j to a constant value in Theorem 1 does not always maximize the equivocation rates. We rewrite the bounds on the equivocation rates in Theorem 1 as follows:

$$\Delta_1 \leq h(X|E) - I(U_1; Y_1|U_2) + I(U_1; Y_1|X) + I(V_1; E) - I(V_1; U_2), \quad (155)$$

$$\Delta_2 \leq h(X|E) - I(U_2; Y_2|U_1) + I(U_2; Y_2|X) + I(V_2; E) - I(V_2; U_1), \quad (156)$$

$$\begin{aligned} \Delta_1 + \Delta_2 &\leq 2h(X|E) - I(U_1, U_2; Y_1, Y_2) \\ &\quad + I(U_1; Y_1|X) + I(U_2; Y_2|X) \\ &\quad + I(V_1; E) + I(V_2; E) - I(V_1; V_2), \end{aligned} \quad (157)$$

$$\begin{aligned} \Delta_1 - R_2 &\leq h(X|E) - I(U_2; Y_2|U_1) - I(U_1; Y_1) \\ &\quad + I(U_1; Y_1|X) + I(V_1; E), \end{aligned} \quad (158)$$

$$\begin{aligned} \Delta_2 - R_1 &\leq h(X|E) - I(U_1; Y_1|U_2) - I(U_2; Y_2) \\ &\quad + I(U_2; Y_2|X) + I(V_2; E). \end{aligned} \quad (159)$$

Then, considering the Markov chains $V_j - X - (E, U_{j'})$ for $j' \neq j \in \{1, 2\}$ and the fact that the variables X and E as well as the auxiliary random variables U_1 and U_2 in (135) are Gaussian, we have four possibilities:

- E is less noisy than U_2 w.r.t. $X \Rightarrow I(V_1; E) \geq I(V_1; U_2)$
⇒ setting $V_1 = U_1$ maximizes (155);
- U_2 is less noisy than E w.r.t. $X \Rightarrow I(V_1; E) \leq I(V_1; U_2)$
⇒ setting $V_1 = \emptyset$ maximizes (155);
- E is less noisy than U_1 w.r.t. $X \Rightarrow I(V_2; E) \geq I(V_2; U_1)$
⇒ setting $V_2 = U_2$ maximizes (156);
- U_1 is less noisy than E w.r.t. $X \Rightarrow I(V_2; E) \leq I(V_2; U_1)$
⇒ setting $V_2 = \emptyset$ maximizes (156).

Therefore, we define

$$\begin{aligned} T_1 &:= \max_{V_1} I(V_1; E) - I(V_1; U_2) \\ &= \max \{0, I(U_1; E) - I(U_1; U_2)\}, \end{aligned} \quad (160)$$

$$\begin{aligned} T_2 &:= \max_{V_2} I(V_2; E) - I(V_2; U_1) \\ &= \max \{0, I(U_2; E) - I(U_2; U_1)\}, \end{aligned} \quad (161)$$

and rewrite (155)–(159) as

$$\Delta_1 \leq h(X|E) - I(U_1; Y_1|U_2) + I(U_1; Y_1|X) + T_1, \quad (162)$$

$$\Delta_2 \leq h(X|E) - I(U_2; Y_2|U_1) + I(U_2; Y_2|X) + T_2, \quad (163)$$

$$\begin{aligned} \Delta_1 + \Delta_2 &\leq 2h(X|E) - I(U_1, U_2; Y_1, Y_2) \\ &\quad + I(U_1; Y_1|X) + I(U_2; Y_2|X) \\ &\quad + T_1 + T_2 + I(U_1; U_2) \mathbb{1}_{\mathbb{R}_{>0}}(T_1 + T_2), \end{aligned} \quad (164)$$

$$\begin{aligned} \Delta_1 - R_2 &\leq h(X|E) - I(U_2; Y_2|U_1) - I(U_1; Y_1) \\ &\quad + I(U_1; Y_1|X) + T_1 + I(U_1; U_2) \mathbb{1}_{\mathbb{R}_{>0}}(T_1), \end{aligned} \quad (165)$$

$$\Delta_2 - R_1 \leq h(X|E) - I(U_1; Y_1|U_2) - I(U_2; Y_2)$$

$$+ I(U_2; Y_2|X) + T_2 + I(U_1; U_2) \mathbb{1}_{\mathbb{R}_{>0}}(T_2), \quad (166)$$

where $\mathbb{1}_{\mathbb{R}_{>0}}(\cdot)$ is the indicator function. Based on the definitions of the auxiliary random variables in (135) and r_j in (141), with straightforward calculations we obtain

$$T_1 = \max \left\{ 0, \frac{1}{2} \log \left(1 + \frac{\frac{\sigma_{N_2}^2}{1-2^{-2r_2}} - \sigma_{N_E}^2}{\frac{\sigma_{N_1}^2}{1-2^{-2r_1}} + \sigma_{N_E}^2} \right) \right\}, \quad (167)$$

$$T_2 = \max \left\{ 0, \frac{1}{2} \log \left(1 + \frac{\frac{\sigma_{N_1}^2}{1-2^{-2r_1}} - \sigma_{N_E}^2}{\frac{\sigma_{N_2}^2}{1-2^{-2r_2}} + \sigma_{N_E}^2} \right) \right\}, \quad (168)$$

$$I(U_1; U_2) = \frac{1}{2} \log \left(\frac{\sigma_X^2 + \frac{\sigma_{N_1}^2}{1-2^{-2r_1}}}{\frac{\sigma_{N_1}^2}{1-2^{-2r_1}} + \frac{\sigma_{N_2}^2}{1-2^{-2r_2}}} \right). \quad (169)$$

Using (162)–(169) and the proof given in Appendix F-A, inequalities of Theorem 4 follows. ■

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